

SparseP

Towards Efficient Sparse Matrix Vector Multiplication
on Real Processing-In-Memory Architectures

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UNIVERSIDAD
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Our Work

Efficient Algorithmic Designs

The first open-source Sparse Matrix Vector Multiplication (SpMV) software package, **SparseP**, for real Processing-In-Memory (PIM) systems

SparseP is Open-Source

SparseP: <https://github.com/CMU-SAFARI/SparseP>

Extensive Characterization

The first comprehensive analysis of SpMV on the first real commercial PIM architecture

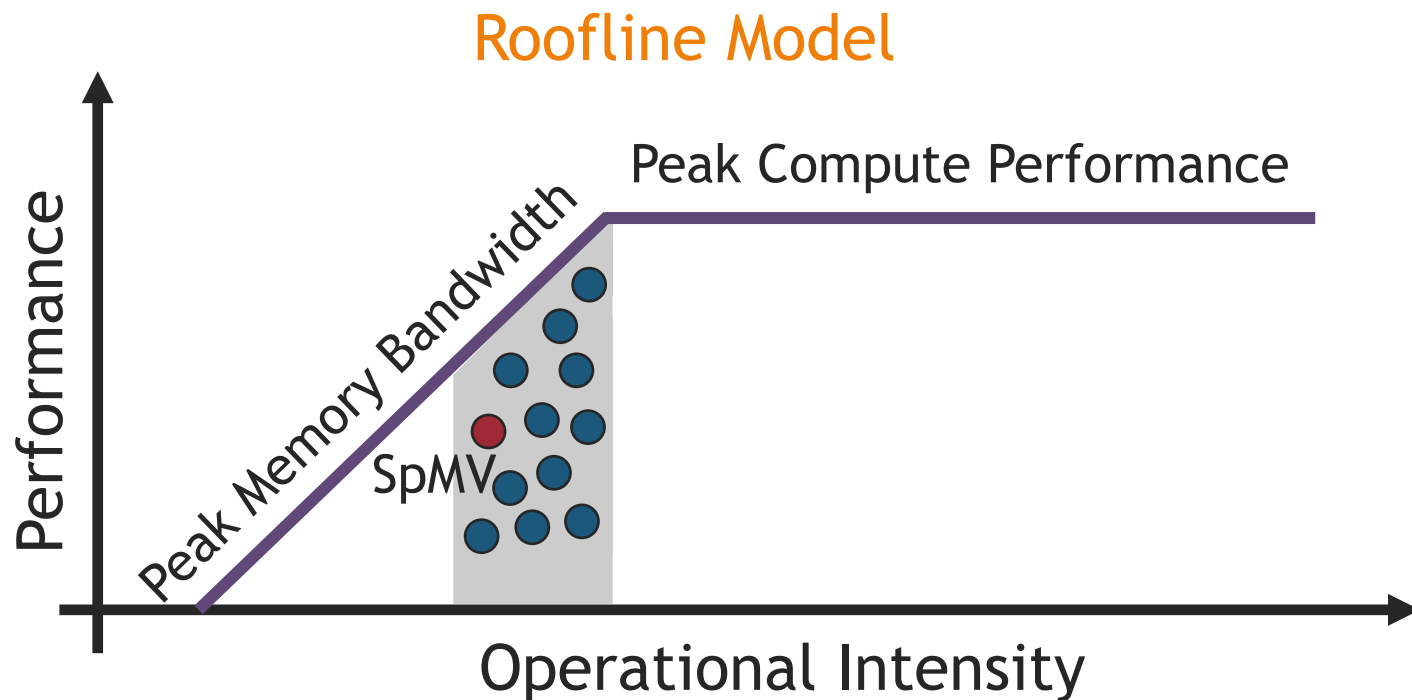
Recommendations for Architects and Programmers

Full Paper: <https://arxiv.org/pdf/2201.05072.pdf>

Sparse Matrix Vector Multiplication

Sparse Matrix Vector Multiplication (SpMV):

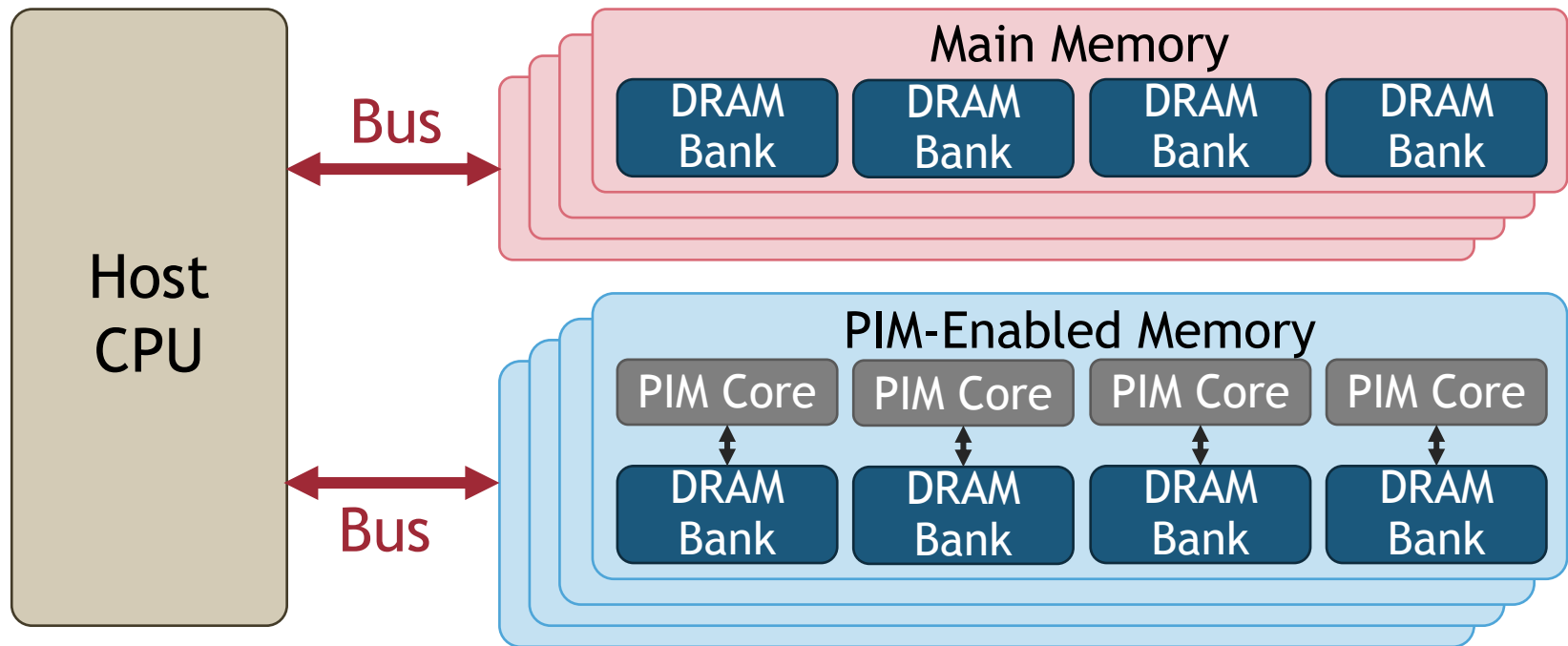
- **Widely-used** kernel in graph processing, machine learning, scientific computing ...
- A **highly memory-bound** kernel



Real Processing-In-Memory Systems

Real **Near-Bank** Processing-In-Memory (**PIM**) Systems:

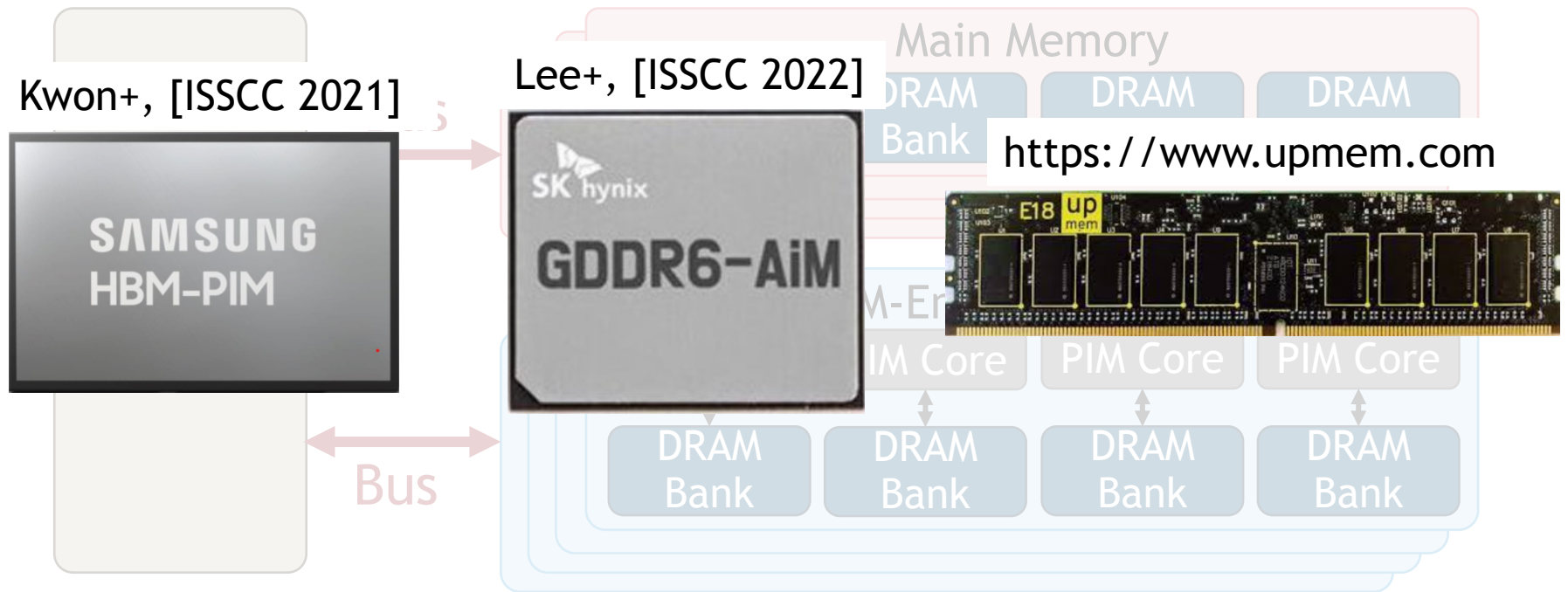
- High levels of parallelism
- Low memory access latency
- Large aggregate memory bandwidth



Real Processing-In-Memory Systems

Real **Near-Bank** Processing-In-Memory (**PIM**) Systems:

- High levels of parallelism
- Low memory access latency
- Large aggregate memory bandwidth



SparseP: SpMV Library for Real PIMs

Our Contributions:

1. Design **efficient SpMV kernels** for current and future PIM systems
 - **25 SpMV kernels**
 - 4 compressed matrix formats (CSR, COO, BCSR, BCOO)
 - 6 data types
 - 4 data partitioning techniques
 - Various load balancing schemes among PIM cores/threads
 - 3 synchronization approaches
2. Provide a **comprehensive analysis** of SpMV on the first commercially-available **real PIM system** **up mem**
 - **26** sparse matrices
 - Comparisons to state-of-the-art **CPU** and **GPU** systems
 - **Recommendations** for software, system and hardware designers

Outline

SpMV Kernels for Real PIM Systems

Key Takeaways from Our Study

Conclusion

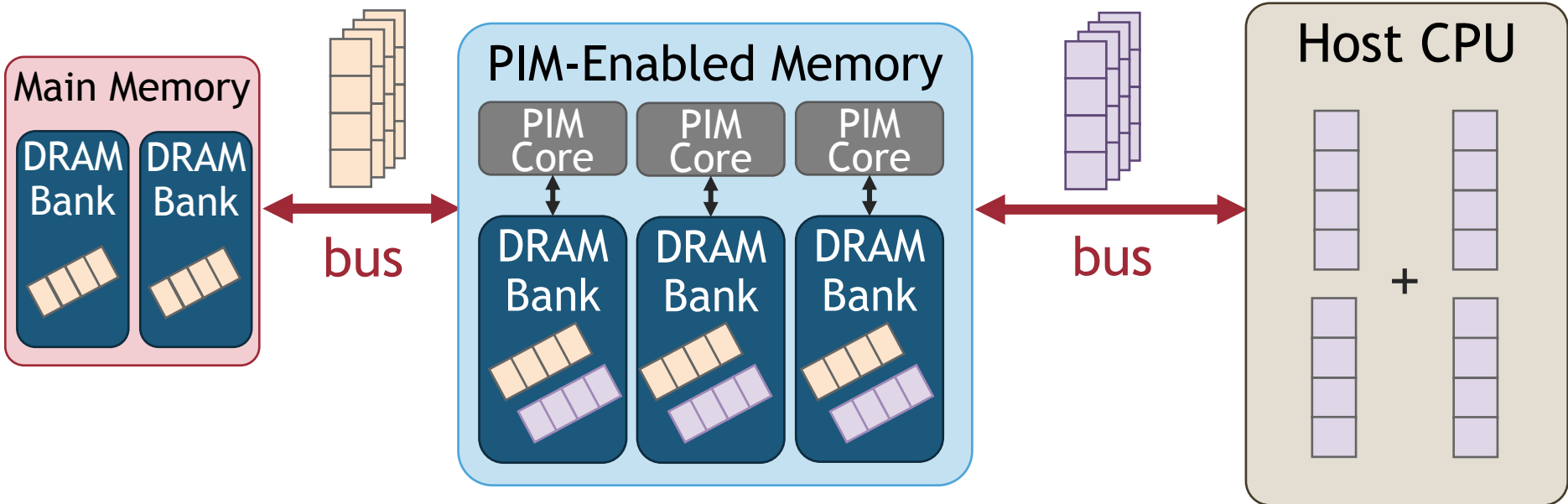
SpMV Execution on a PIM System

①
Load the
input vector

②
Execute the
kernel

③
Retrieve the
partial results

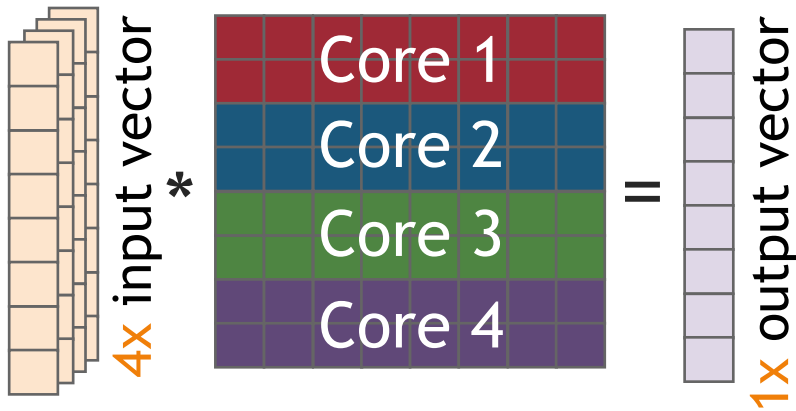
④
Merge the
partial results



Data Partitioning Techniques

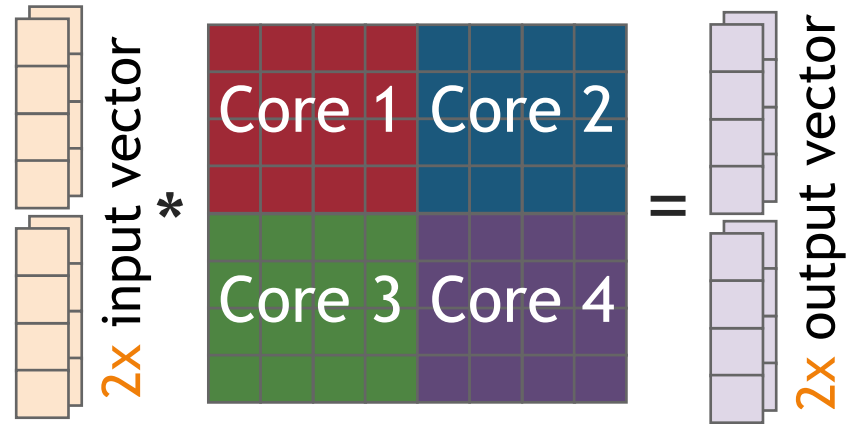
SparseP supports two types of data partitioning techniques:

1D Partitioning



perform the **complete**
SpMV computation
only on PIM cores

2D Partitioning



trade-off
computation vs
data transfer costs

1D Partitioning Technique

Load-Balancing Approaches:

- CSR, COO:
 - Balance Rows
 - Balance NNZs *
- BCSR, BCOO:
 - Balance Blocks ^
 - Balance NNZs ^

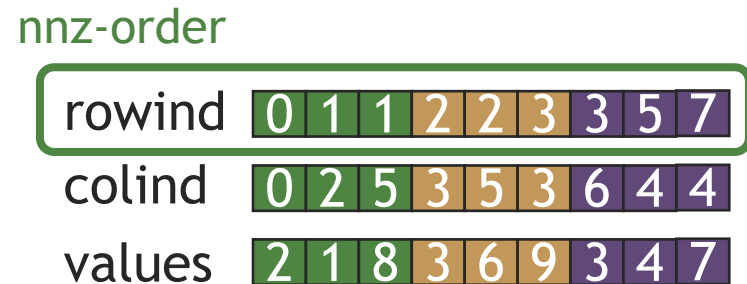
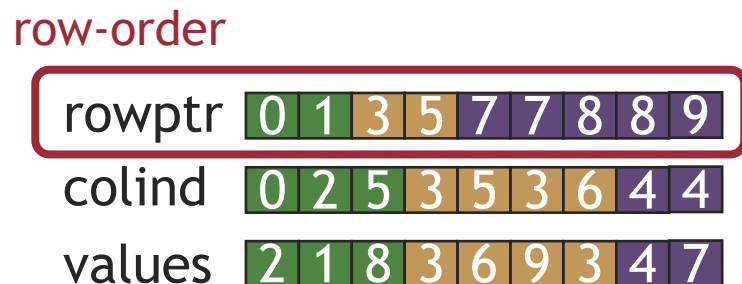
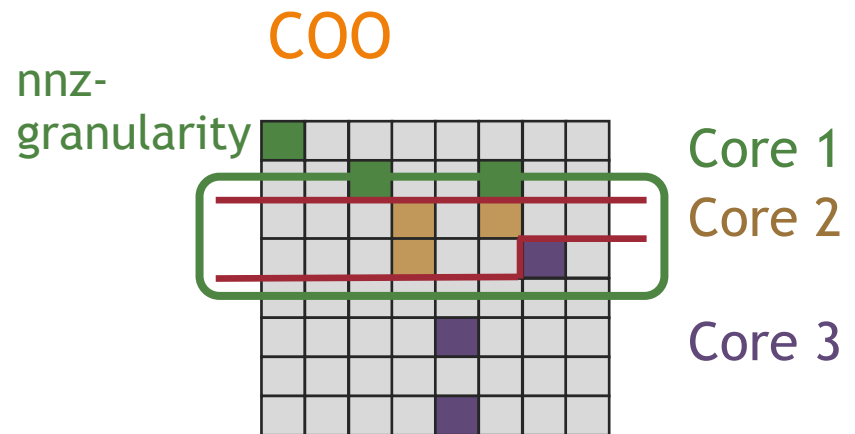
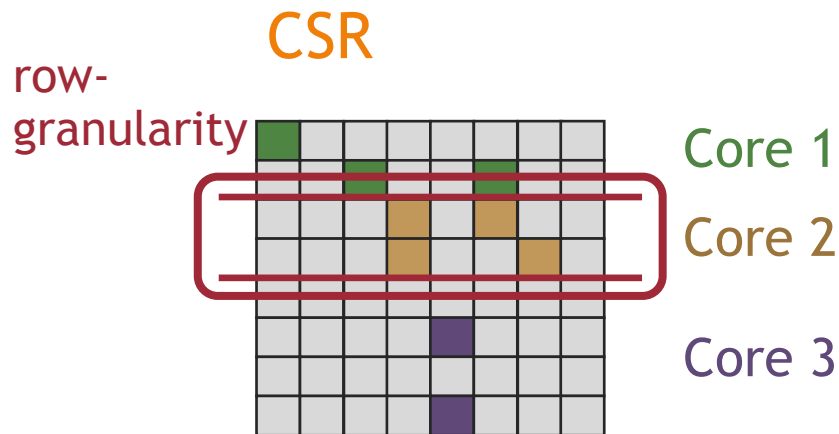
* row-granularity for CSR

^ block-row-granularity for BCSR

1D Partitioning Technique

Load-Balancing of #NNZs:

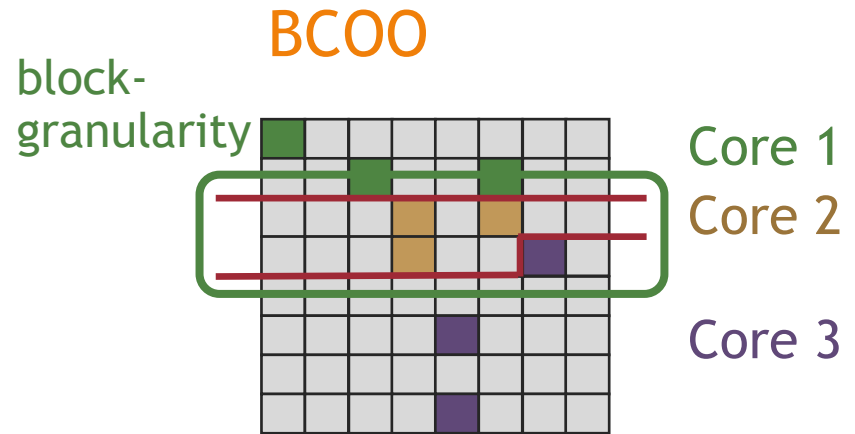
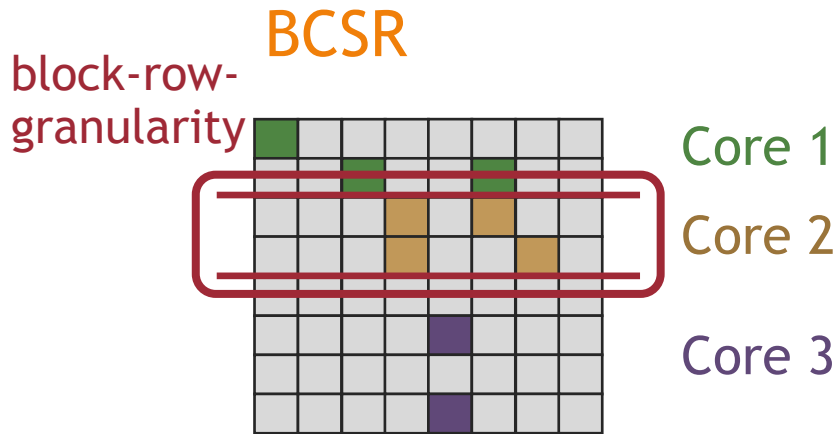
- CSR (row-granularity), COO



1D Partitioning Technique

Load-Balancing of #NNZs:

- CSR (row-granularity), COO
- BCSR (block-row-granularity), BCOO



block-row-order

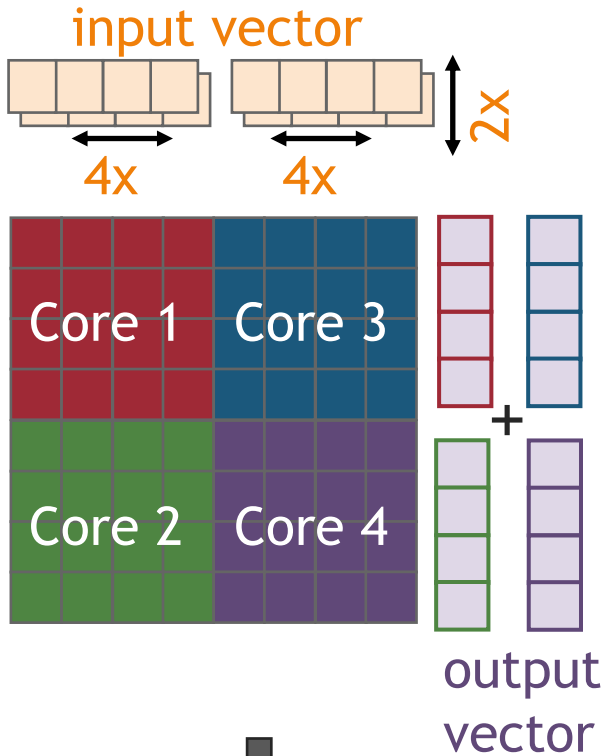
rowptr	0	1	3	5	7	7	8	8	9
colind	0	2	5	3	5	3	6	4	4
values	2	1	8	3	6	9	3	4	7

block-order

rowind	0	1	1	2	2	3	3	5	7
colind	0	2	5	3	5	3	6	4	4
values	2	1	8	3	6	9	3	4	7

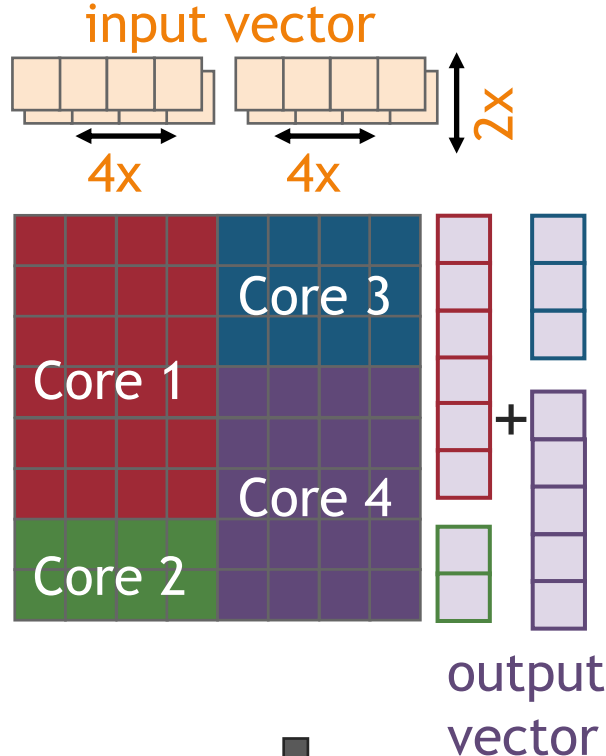
2D Partitioning Technique

Equally-Sized Tiles



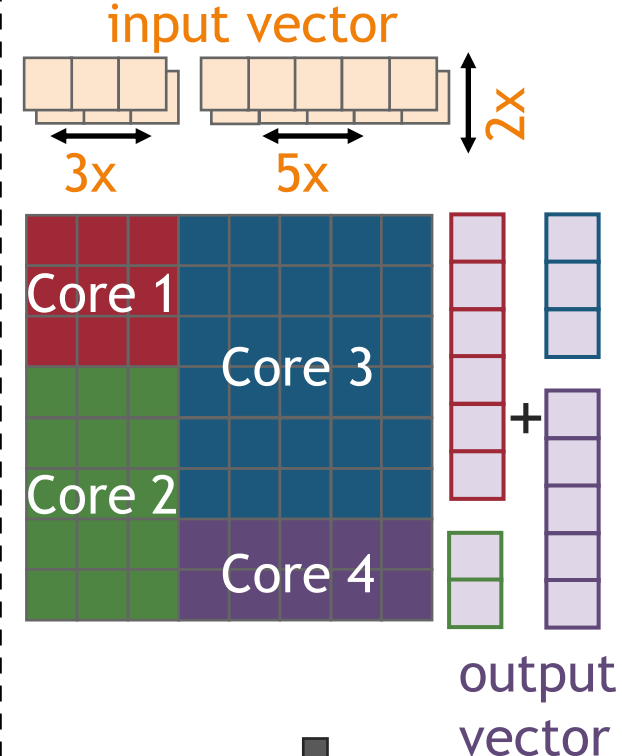
High NNZ **imbalance**
across PIM cores

Equally-Wide Tiles



High NNZ **balance**
across PIM cores of the
same **vertical** partition

Variable-Sized Tiles

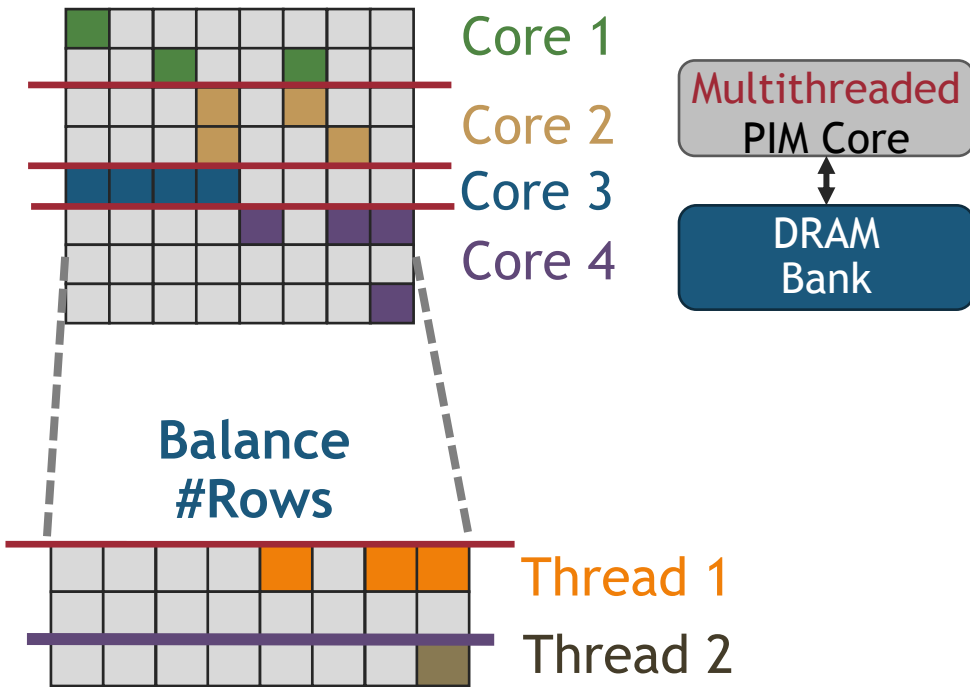


High NNZ **balance**
across **all** PIM cores

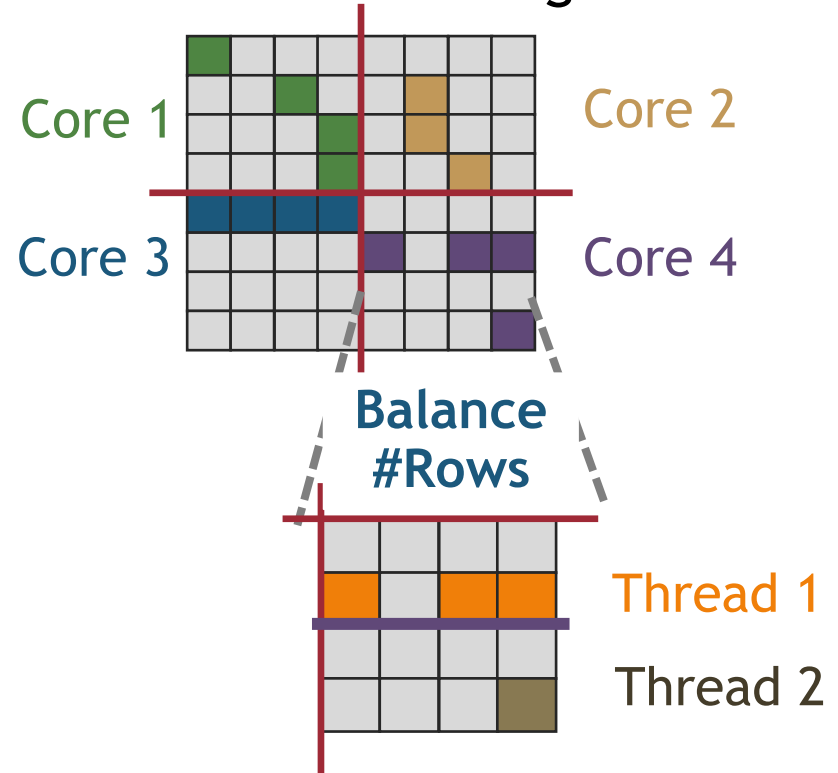
Load-Balance across Threads

Multithreaded PIM Cores:

1D Partitioning



2D Partitioning

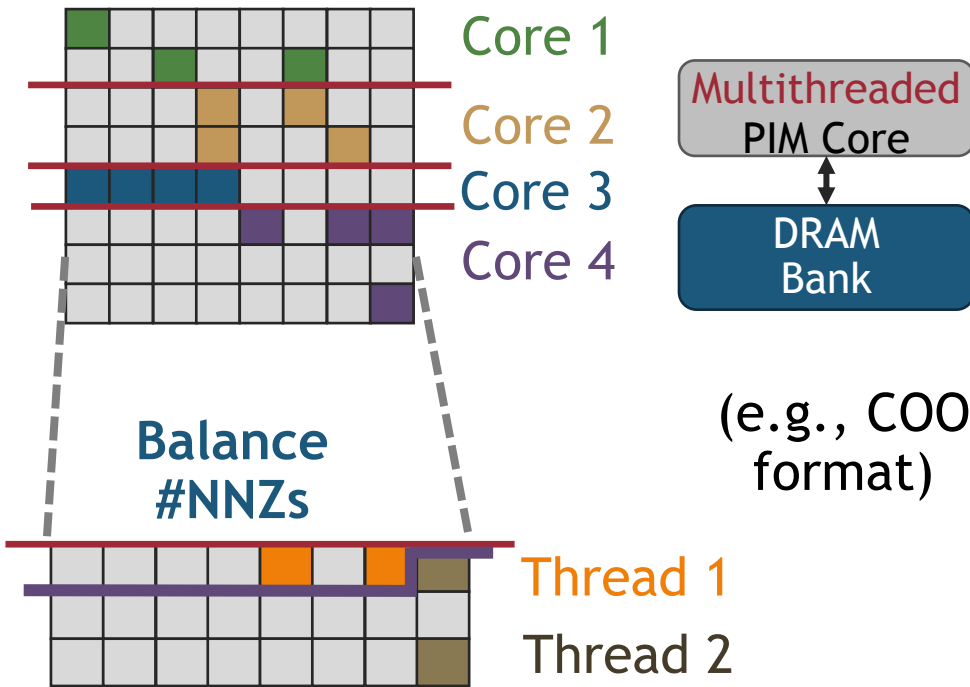


- Various **load-balance** schemes across threads

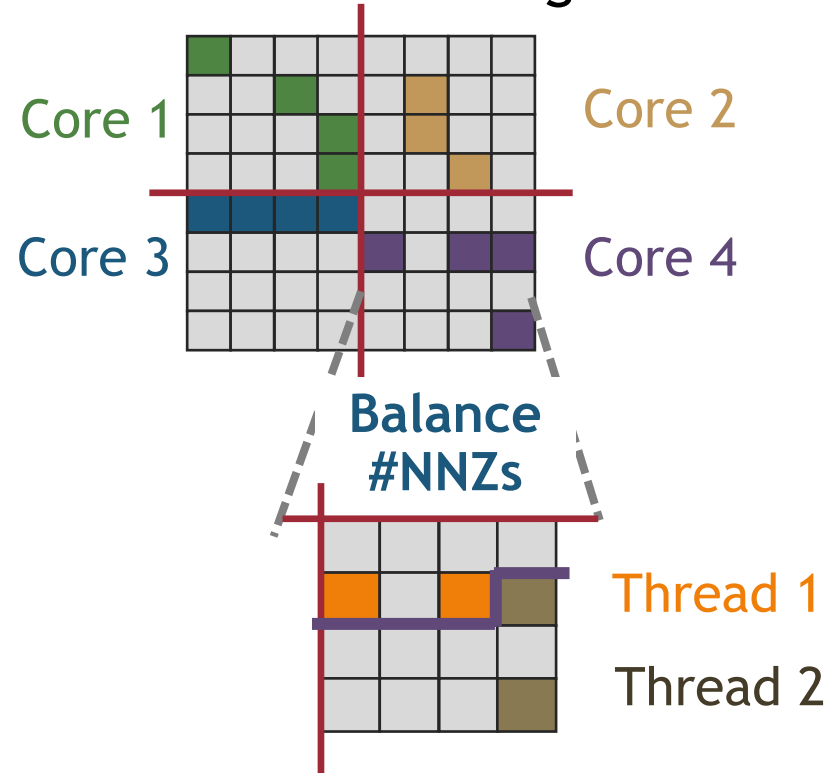
Load-Balance across Threads

Multithreaded PIM Cores:

1D Partitioning



2D Partitioning

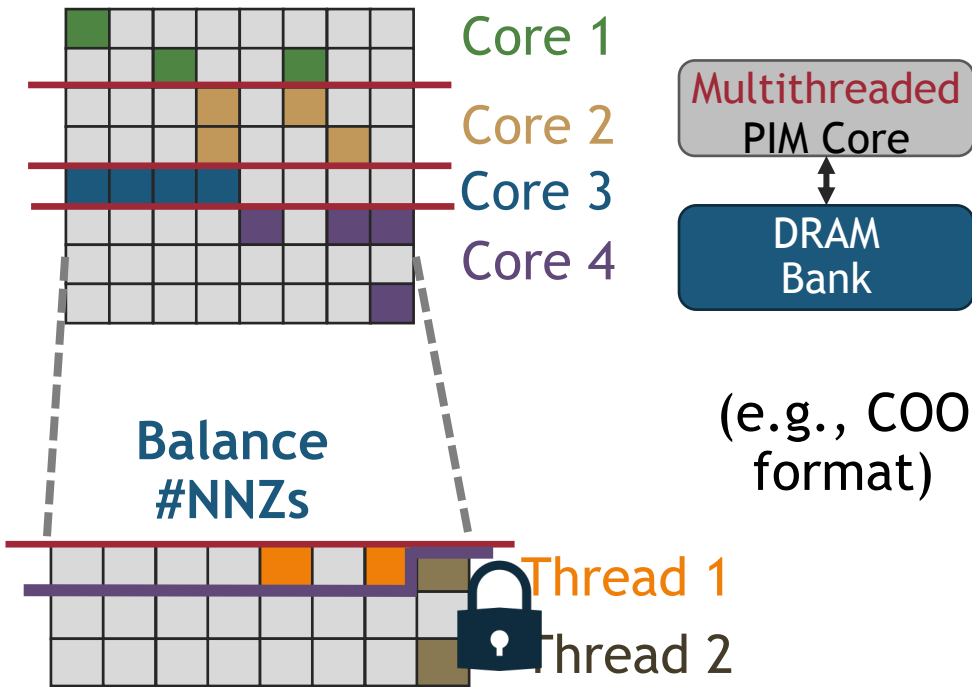


- Various load-balance schemes across threads

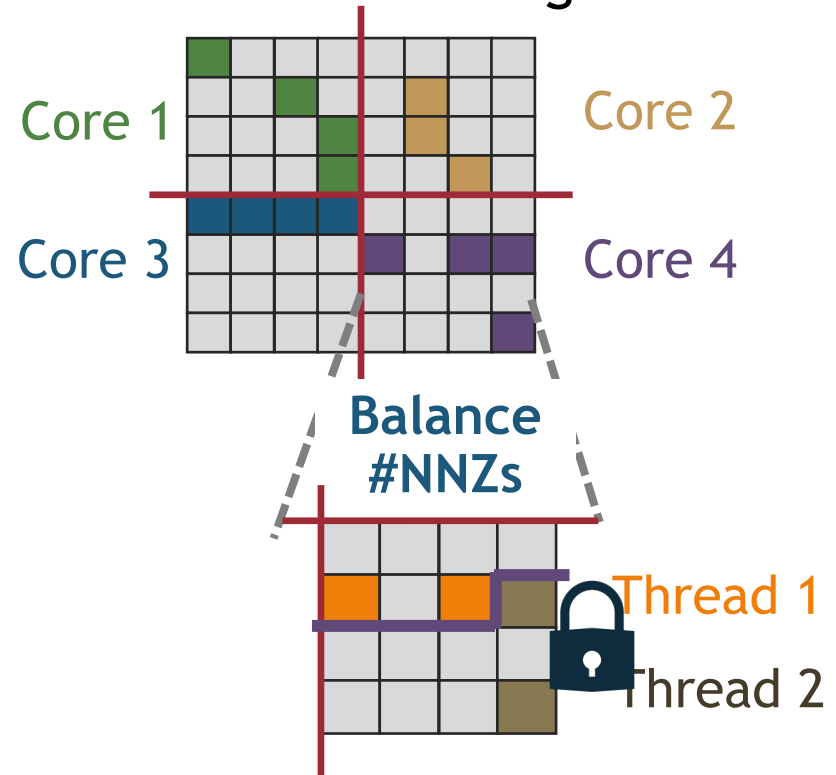
Load-Balance across Threads

Multithreaded PIM Cores:

1D Partitioning

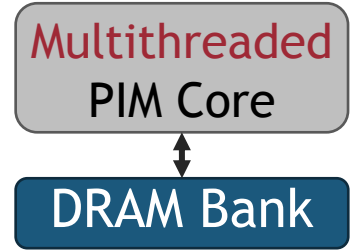


2D Partitioning



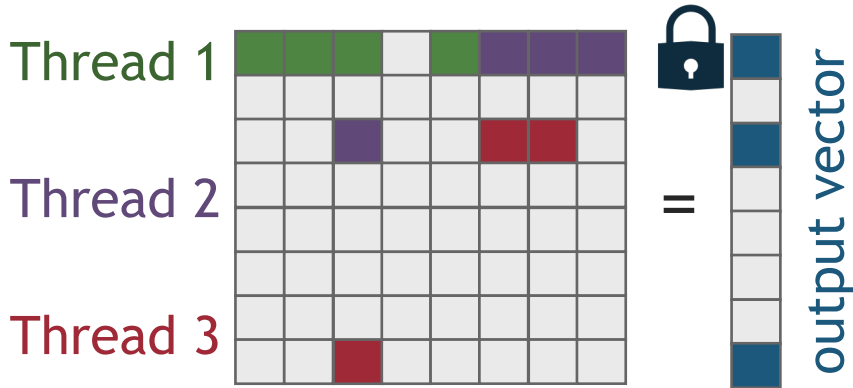
- Various **load-balance** schemes across threads
- Various **synchronization** approaches among threads

Synchronization Approaches

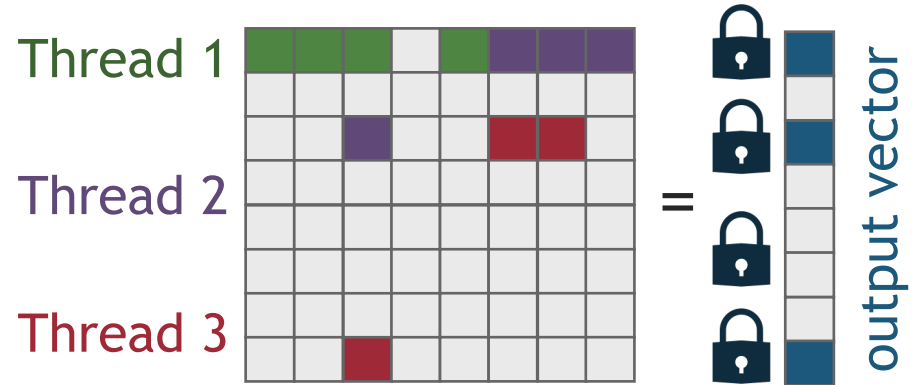


Multithreaded PIM Core:

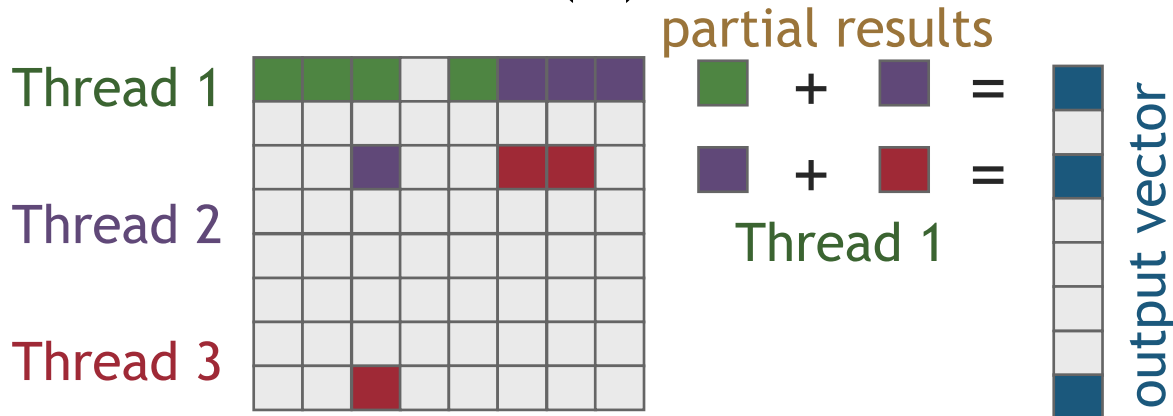
Coarse-Grained (lb-cg)



Fine-Grained (lb-fg)



Lock-Free (lf)



SparseP Software Package

25 SpMV kernels for PIM Systems →

<https://github.com/CMU-SAFARI/SparseP>

Partitioning	Matrix Format	Load-Balancing
9x 1D Kernels	CSR	rows, nnzs *
	COO [▲]	rows, nnzs *, nnzs
	BCSR	blocks [^] , nnzs [^]
	BCOO [▲]	blocks, nnzs
4x 2D Equally-Sized Tiles	CSR	--
	COO [▲]	--
	BCSR	--
	BCOO [▲]	--
6x 2D Equally-Wide Tiles	CSR	nnzs *
	COO [▲]	nnzs
	BCSR	blocks [^] , nnzs [^]
	BCOO [▲]	blocks, nnzs
6x 2D Variable-Sized Tiles	CSR	nnzs *
	COO [▲]	nnzs
	BCSR	blocks [^] , nnzs [^]
	BCOO [▲]	blocks, nnz

Load-balance

across PIM cores/threads:

* row-granularity (CSR)

[^] block-row-granularity (BCSR)

Synchronization

among threads of a PIM core:

[▲] lb-cg, lb-fb, lf (COO, BCOO)

Data Types:

- 8-bit integer
- 16-bit integer
- 32-bit integer
- 64-bit integer
- 32-bit float
- 64-bit float

Outline

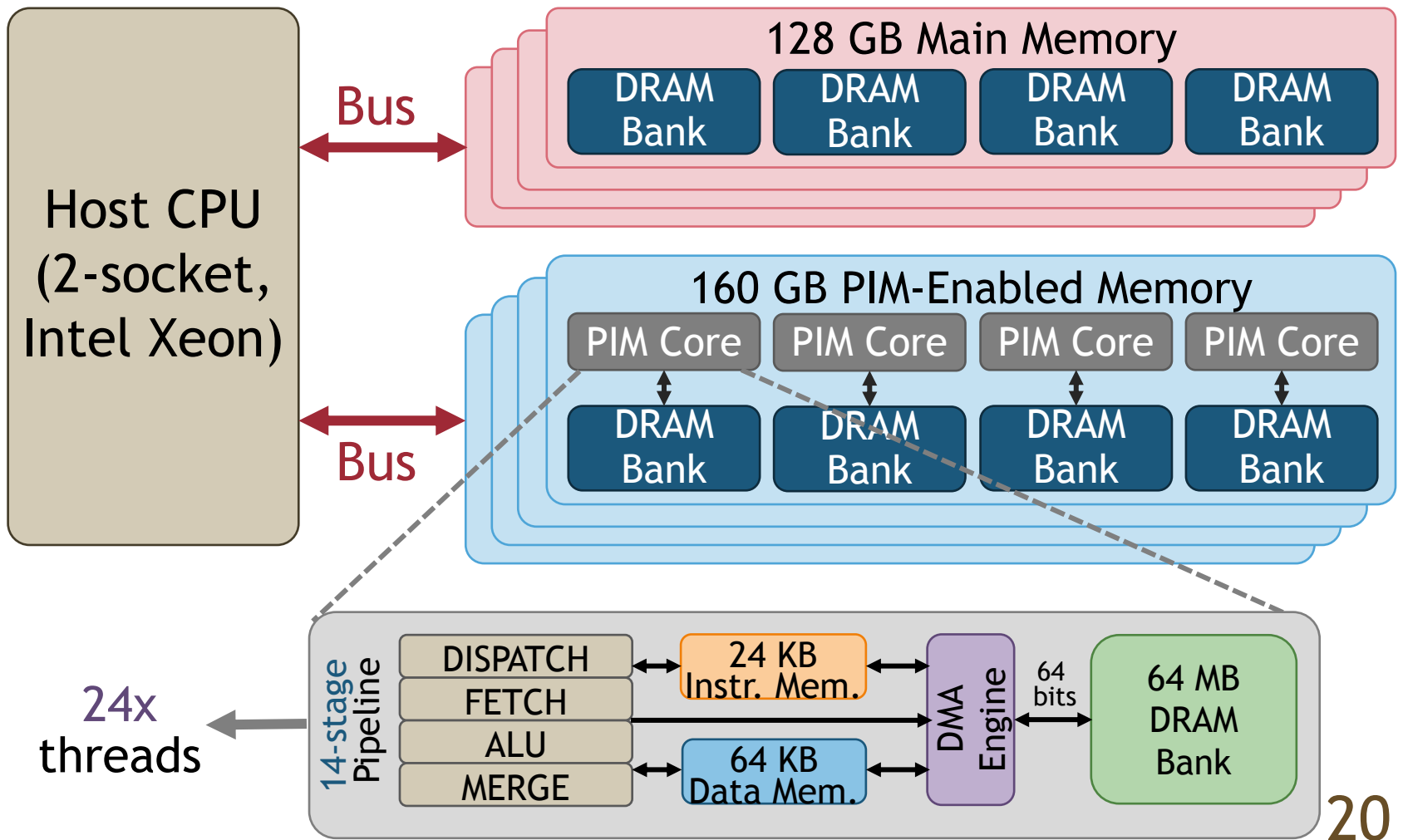
SpMV Kernels for Real PIM Systems

Key Takeaways from Our Study

Conclusion

UPMEM-based PIM System

- 20 UPMEM PIM DIMMs with 2560 PIM cores in total
- Each multithreaded PIM core supports 24 threads

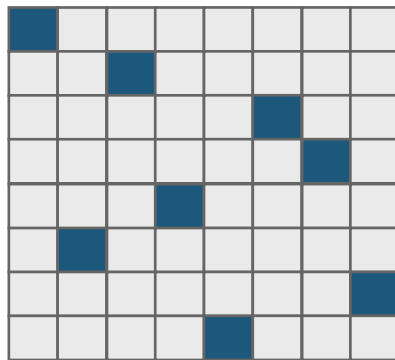


Sparse Matrix Data Set

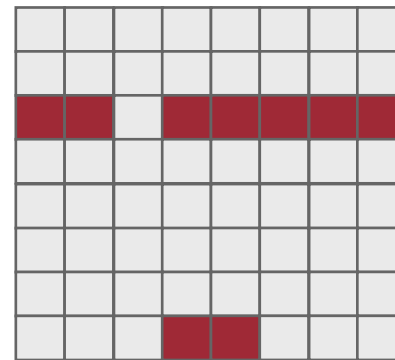
26 sparse matrices*:

- Diverse **sparsity** patterns
- Variability on **irregular** patterns
- Variability on **block** patterns

Regular Matrix



Scale-Free Matrix



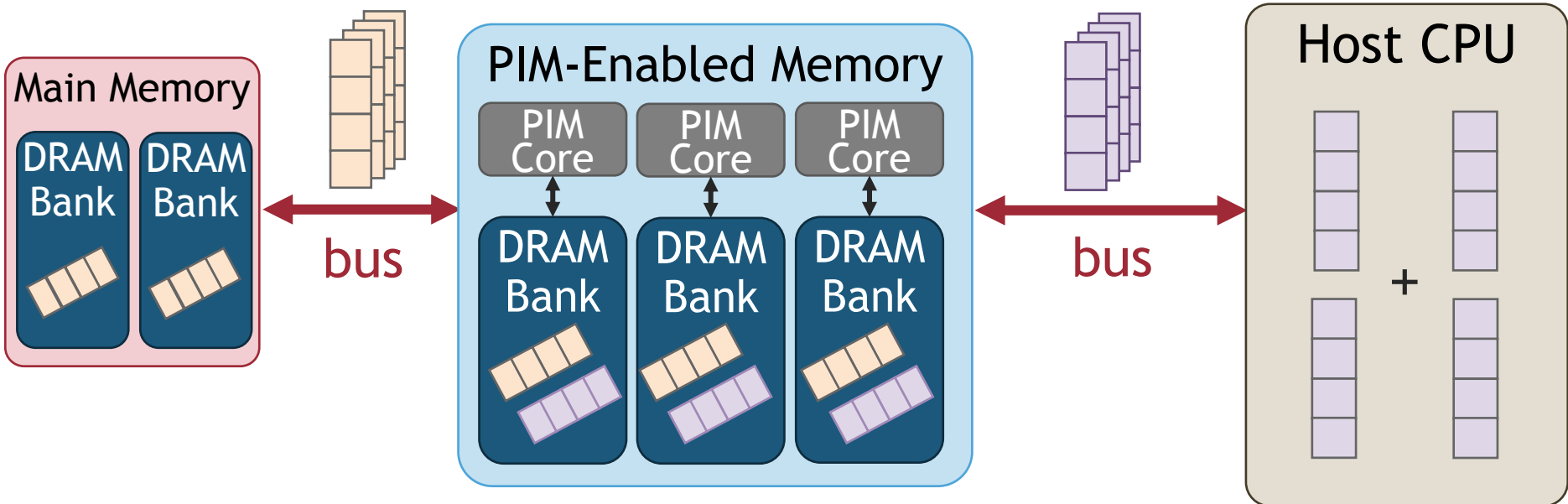
* Suite Sparse Matrix Collection: <https://sparse.tamu.edu/>

Kernel Execution on One PIM Core

① Load the input vector

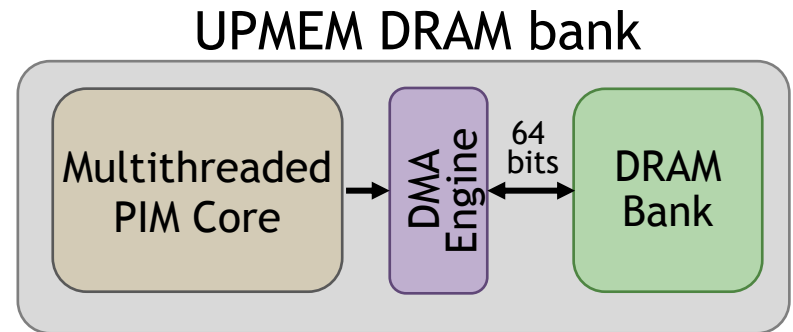
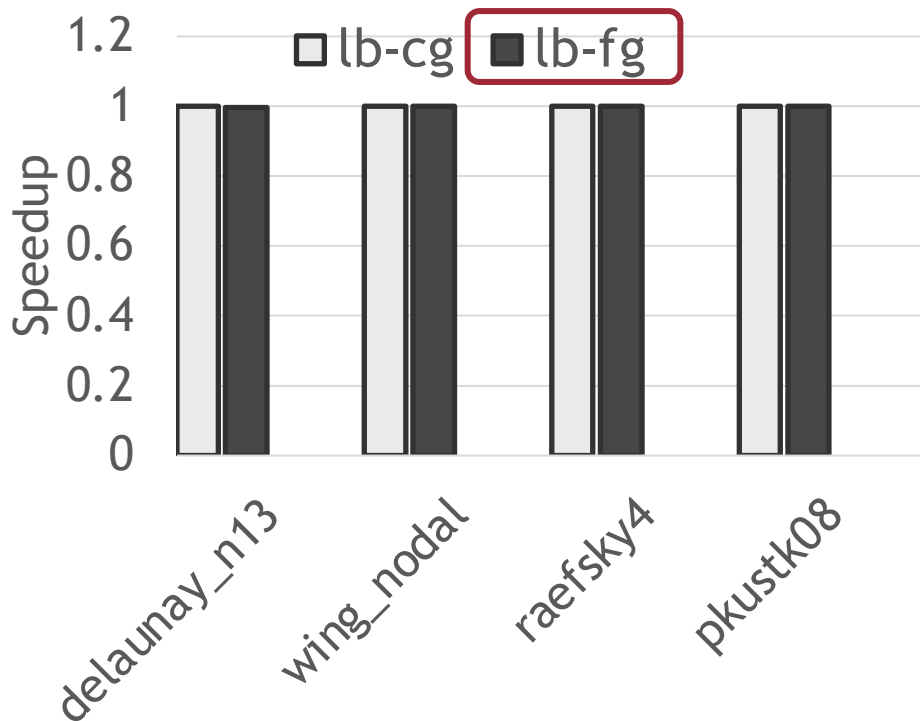
② Execute the kernel

③ Retrieve the partial results
④ Merge the partial results



Lock-Based Synchronization

16 threads, COO, 32-bit integer



Fine-grained locking (lb-fg) does **not** improve performance over coarse-grained locking (lb-cg)

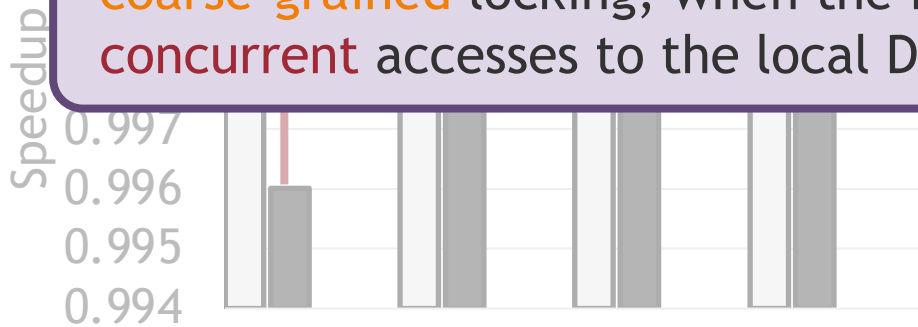
Fine-Grained Locking: memory **accesses** to the **local** DRAM bank are **serialized** in the DMA engine of the UPEM PIM hardware.

Lock-Based Synchronization

16 threads, COO, 32-bit integer

Key Takeaway 1

Fine-grained locking approaches **cannot improve performance** over **coarse-grained** locking, when the PIM hardware does **not** support **concurrent** accesses to the local DRAM bank.



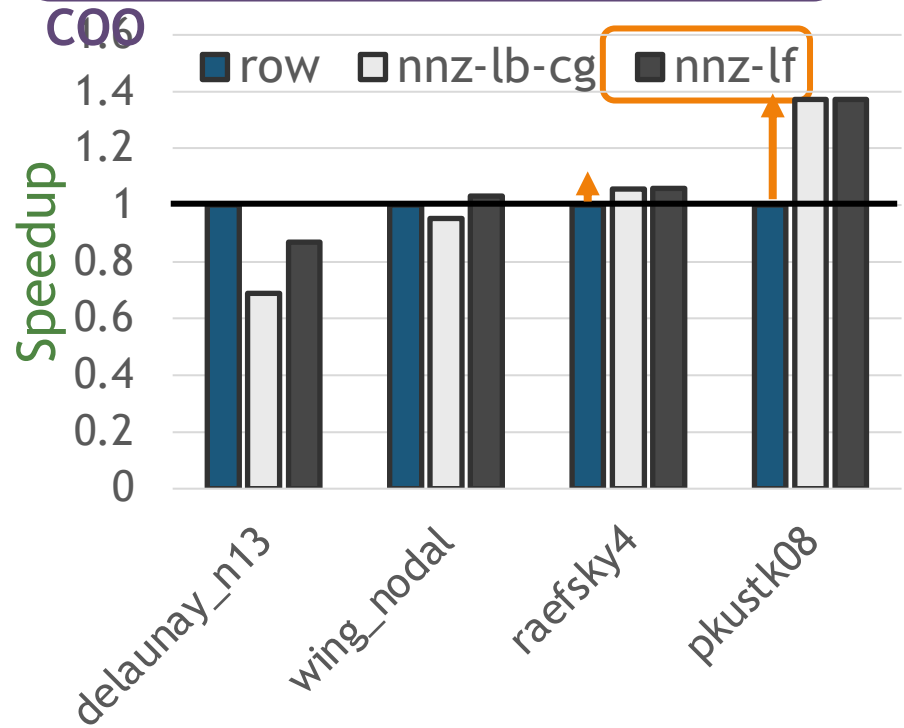
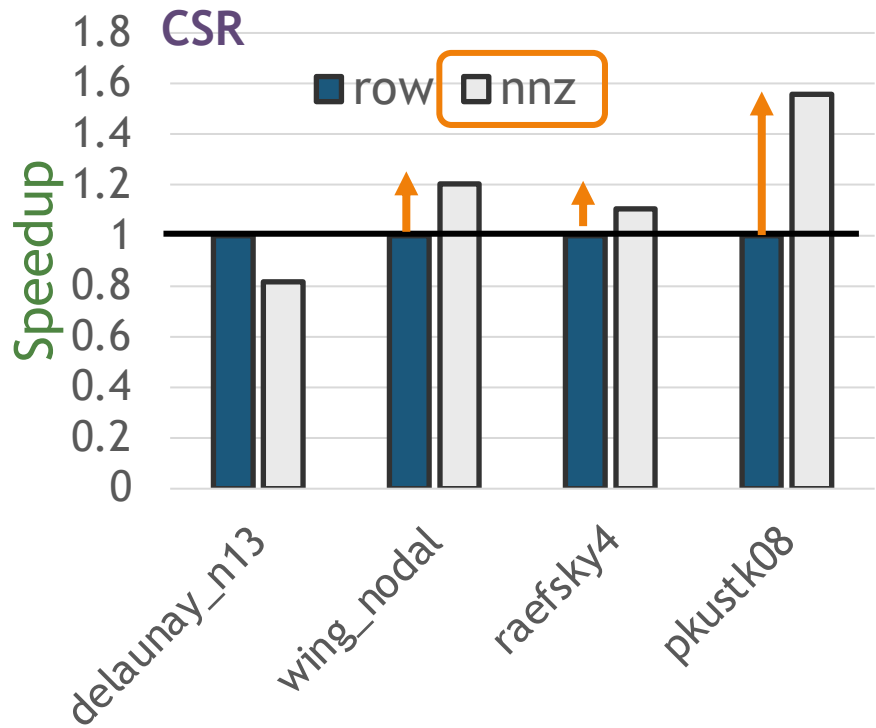
Recommendation 1

Provide **low-cost synchronization** support and hardware support to enable **concurrent** memory **accesses** to the local DRAM bank, and integrate **multiple** DRAM **banks** per PIM core to increase execution parallelism.

Load-Balance within a PIM Core

16 threads, 32-bit integer

Load-balancing #NNZs performs best in most matrices

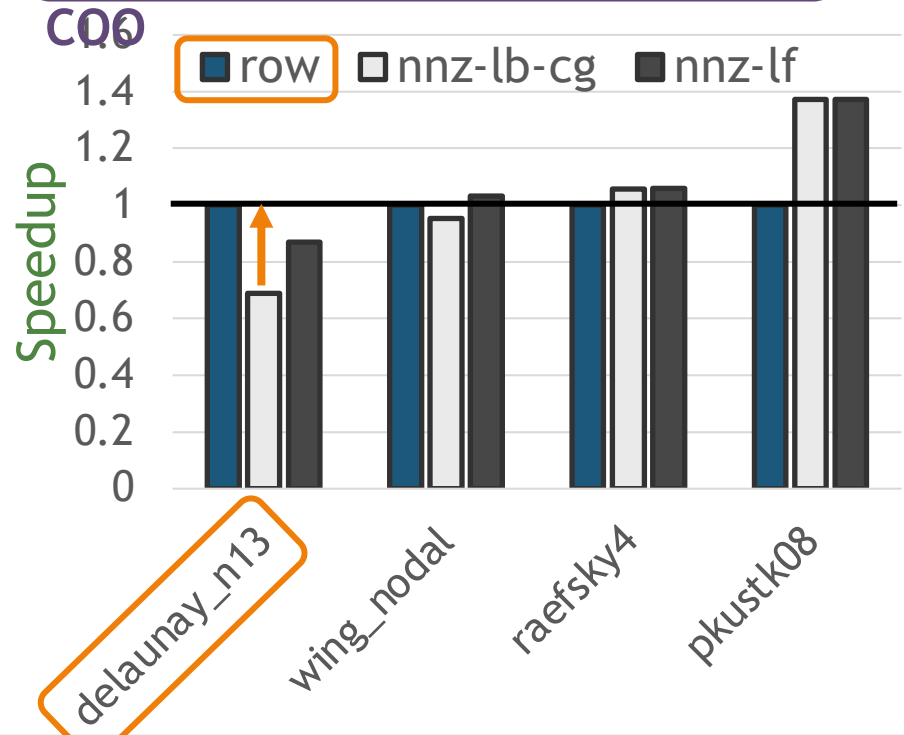
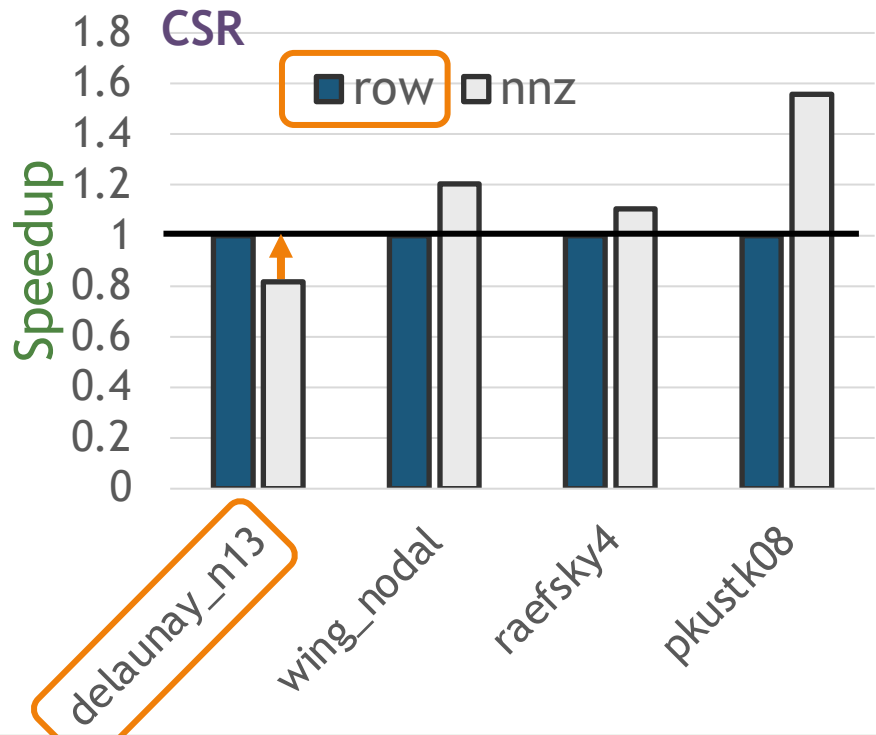


Load-balancing #NNZ typically provides high computation balance across threads of a compute-limited PIM core

Load-Balance within a PIM Core

16 threads, 32-bit integer

Load-balancing #NNZs causes high row imbalance



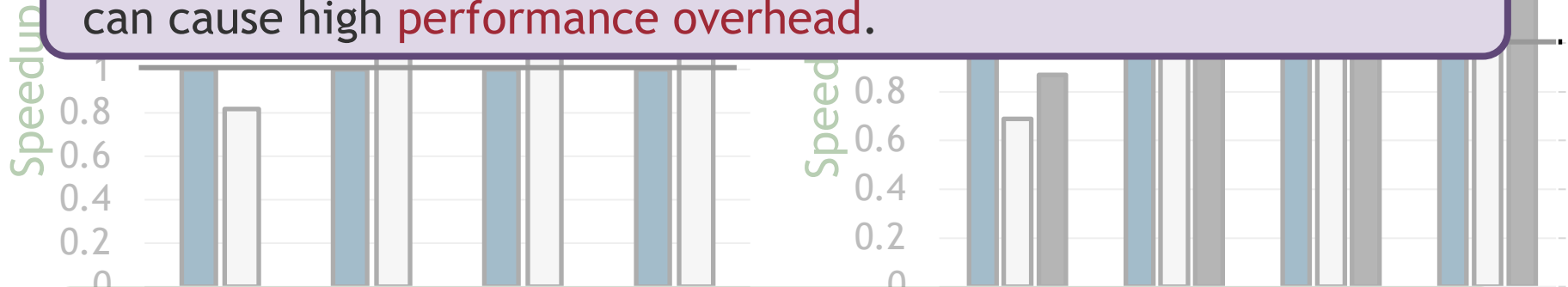
Load-balancing #NNZs: one single thread performs a much higher #memory accesses and #synchronization operations than the rest

Load-Balance within a PIM Core

16 threads, 32-bit integer

Key Takeaway 2

High **operation imbalance** in computation, synchronization, or memory instructions executed by multiple threads of a PIM core can cause high **performance overhead**.

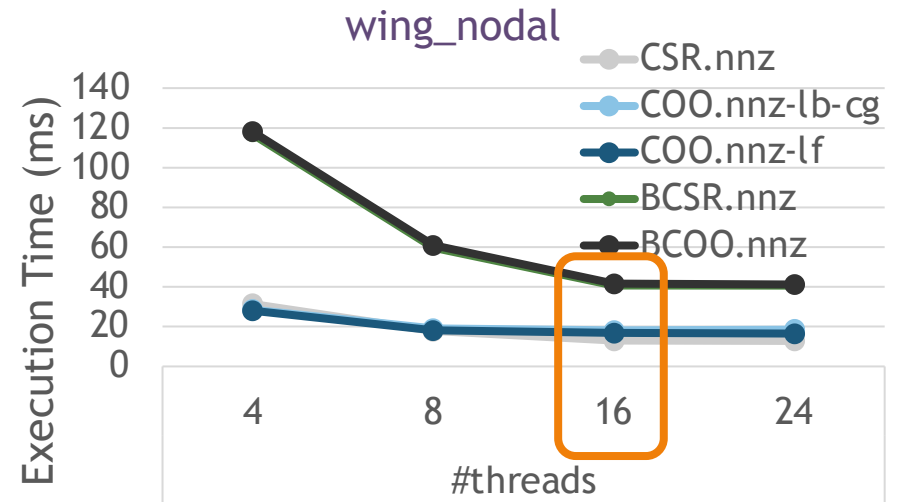
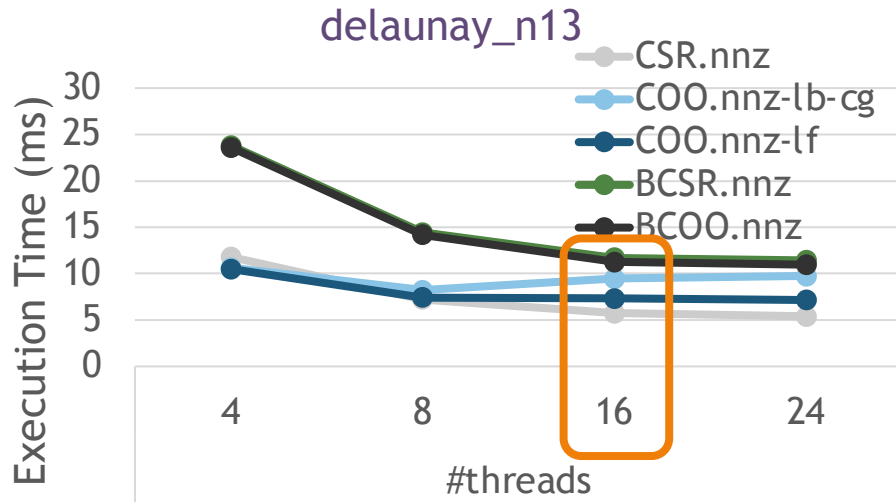


Recommendation 2

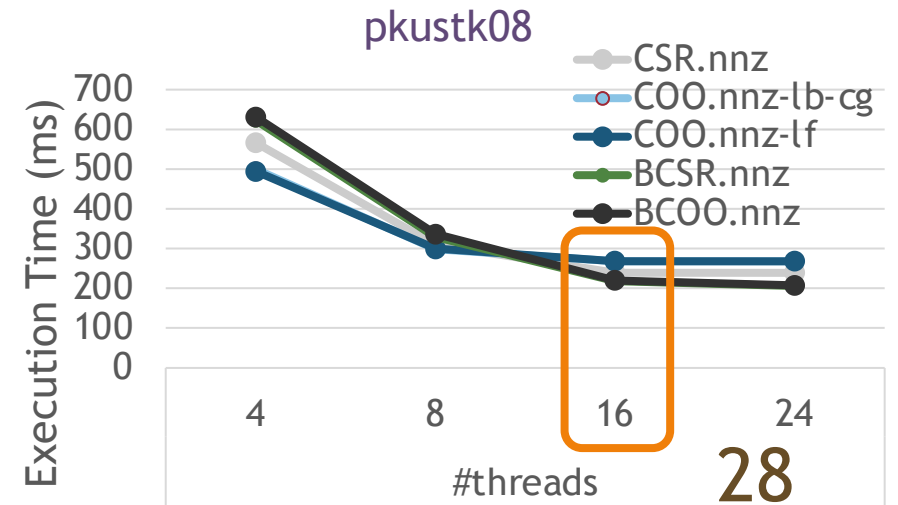
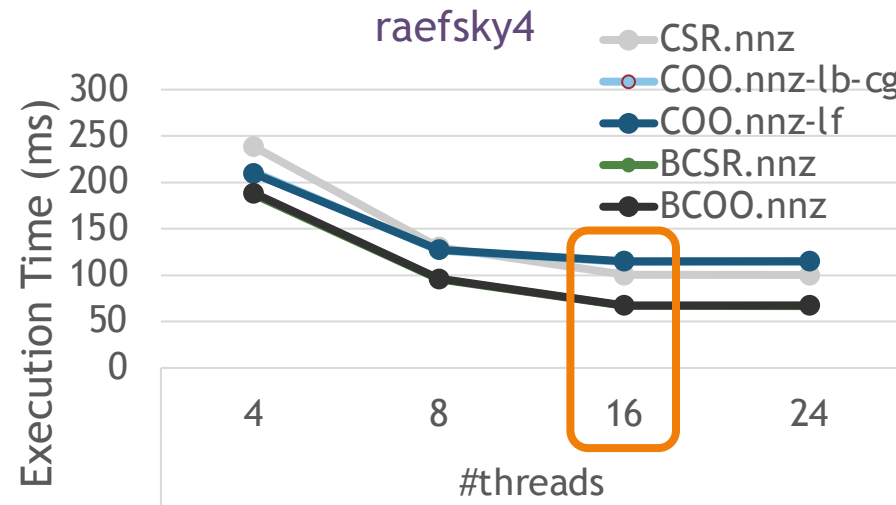
Design algorithms that provide **high load balance** across threads of PIM core in terms of computations, synchronization points and memory accesses.

Scalability within a PIM Core

32-bit integer



Scalability increases up to 16 threads



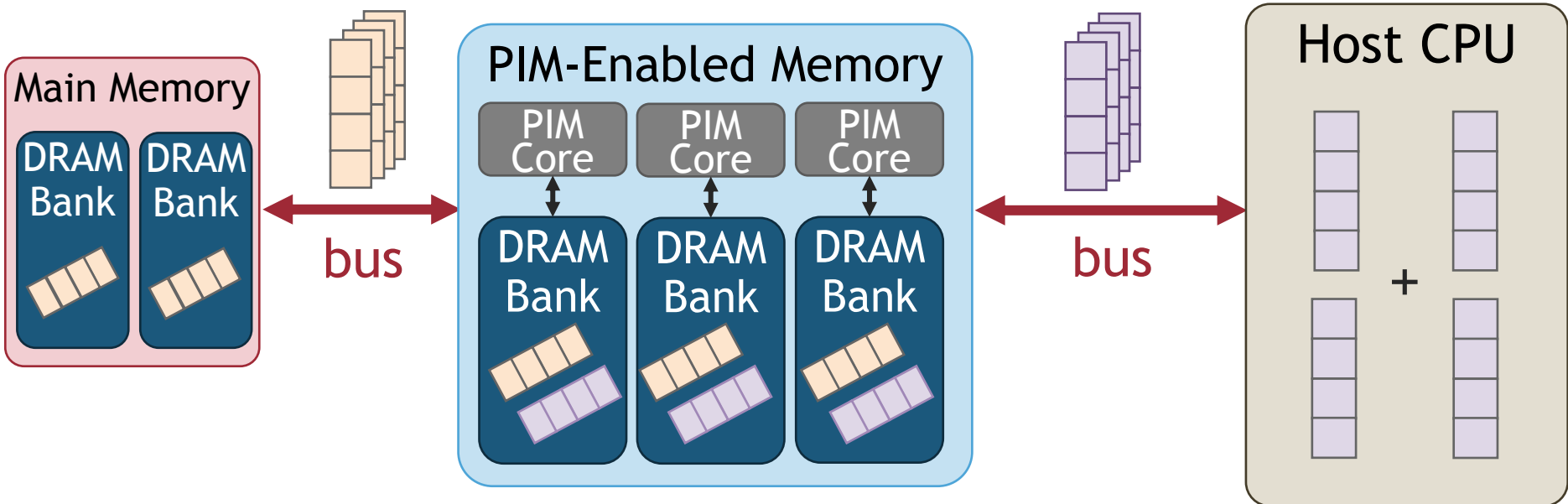
28

Kernel Execution on Multiple PIM Cores

① Load the input vector

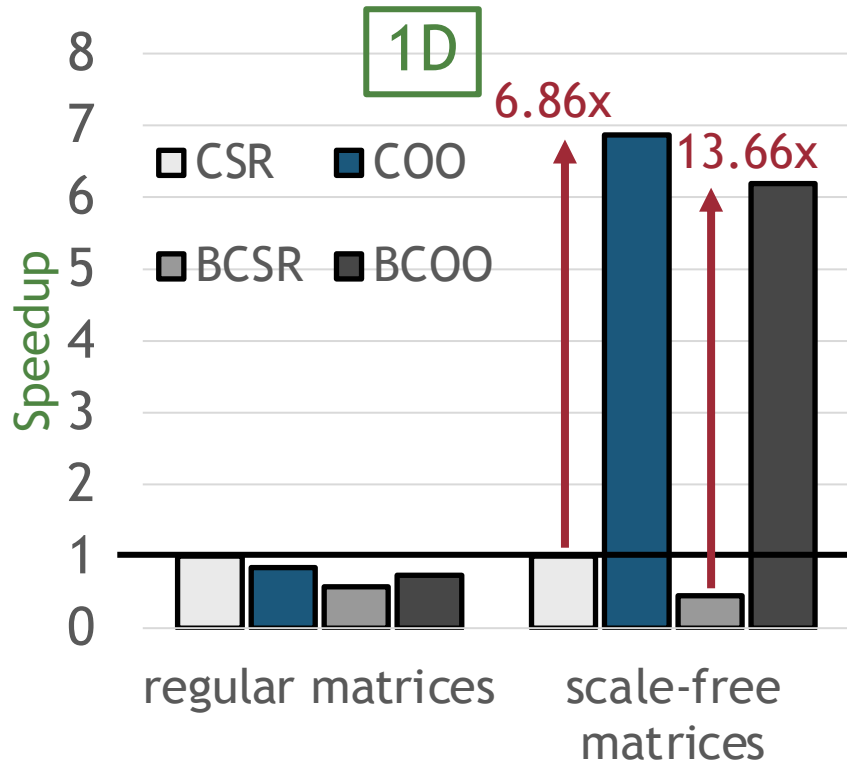
② Execute the kernel

③ Retrieve the partial results
④ Merge the partial results

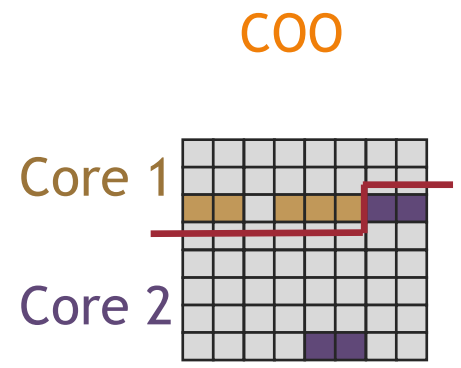
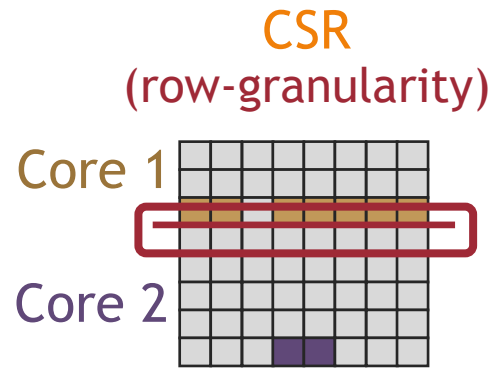


Comparison of Compressed Formats

2048 PIM Cores, 32-bit integer



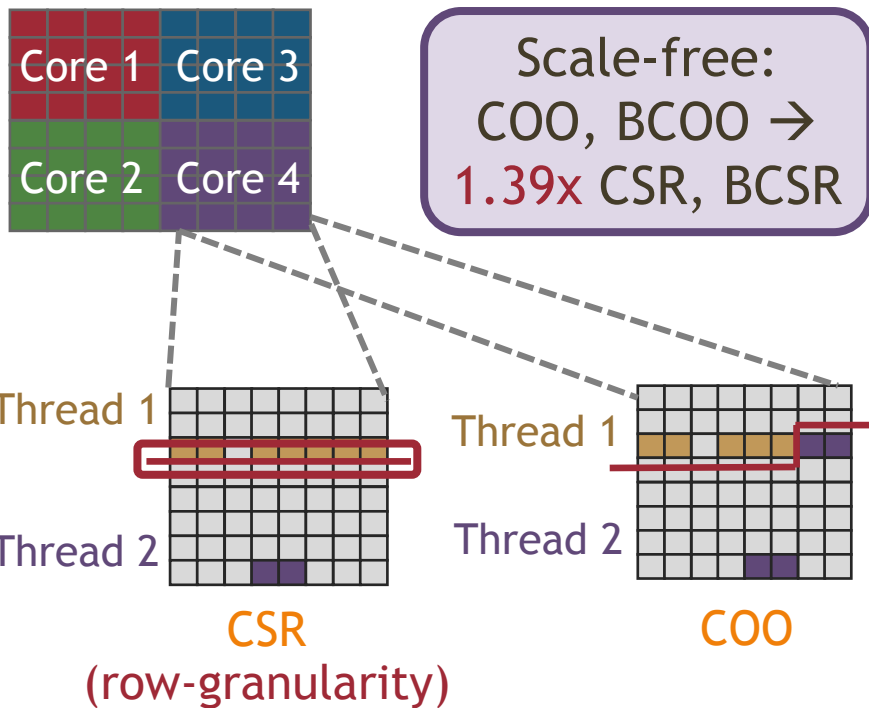
Scale-free: COO, BCOO → 10.26x CSR, BCSR



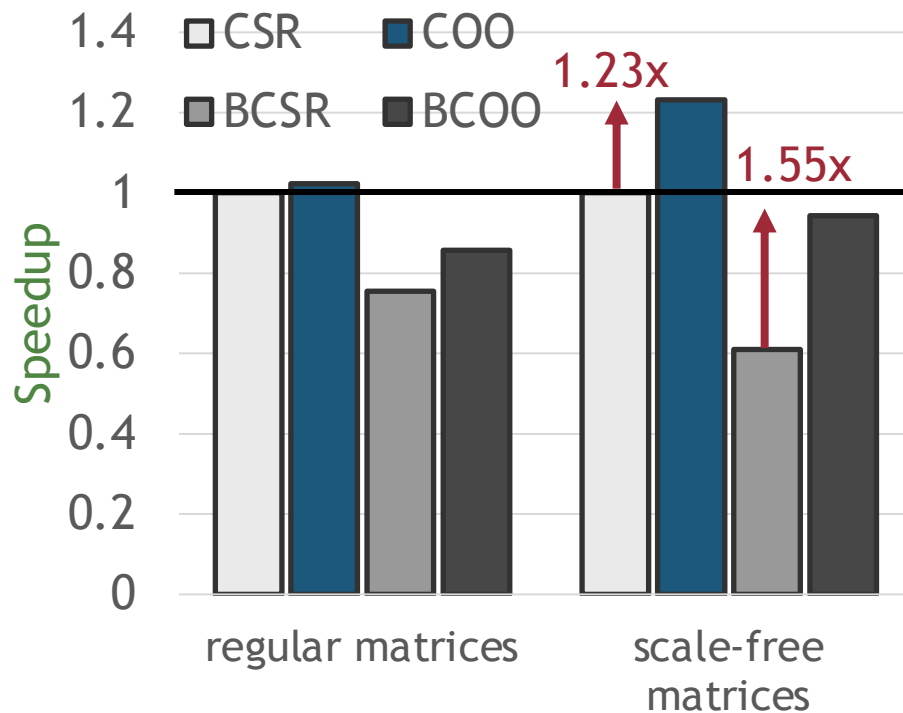
In **scale-free** matrices, **COO + BCOO** provide higher non-zero element balance across PIM cores than **CSR + BCSR**, respectively.

Comparison of Compressed Formats

2048 PIM Cores, 32-bit integer



2D Equally-Sized Tiles



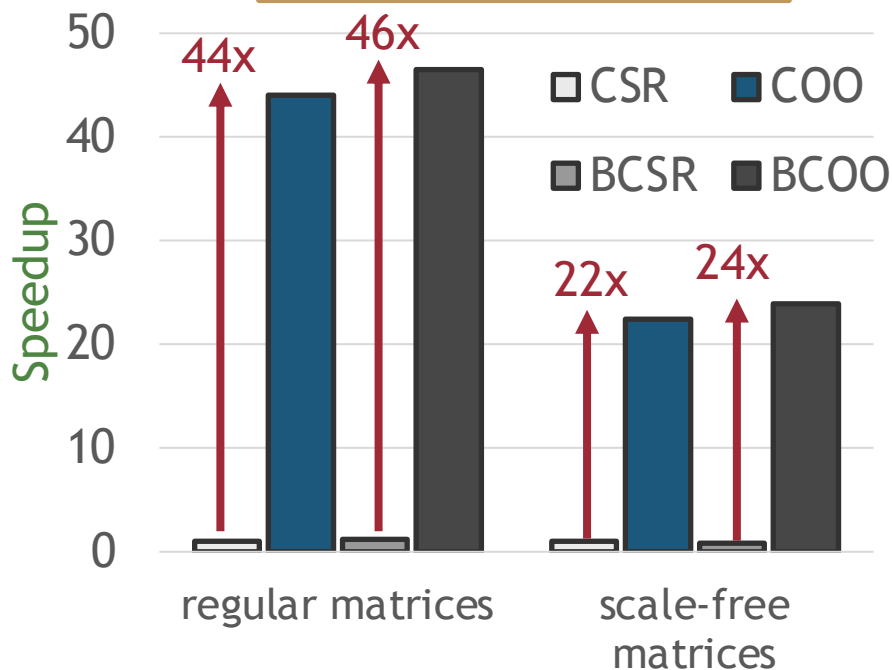
In **scale-free** matrices, **COO** + **BCOO** provide higher non-zero element balance across threads than **CSR** + **BCSR**, respectively.

Comparison of Compressed Formats

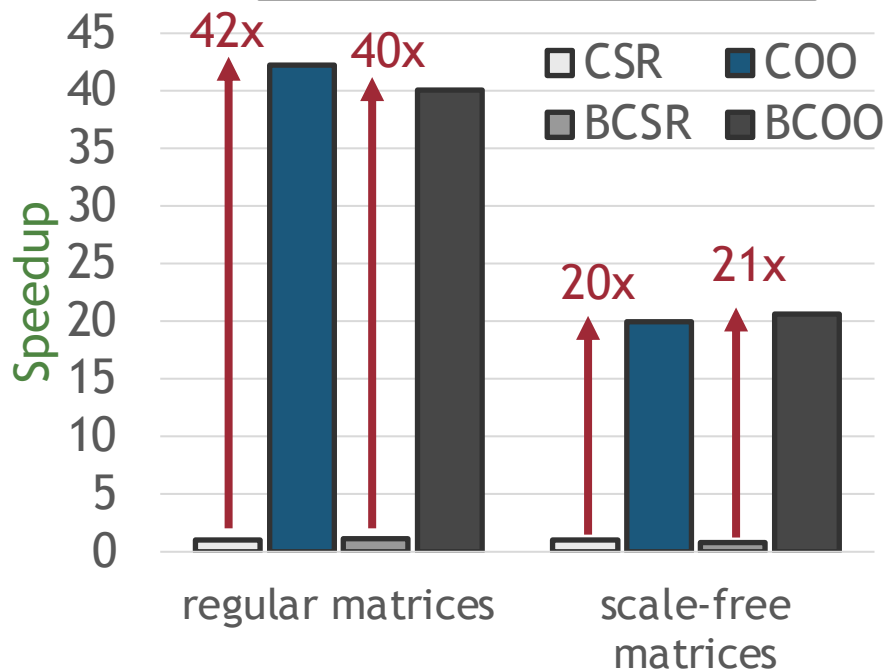
2048 PIM Cores, 32-bit integer

COO, BCOO → 32.38x CSR, BCSR

2D Equally-Wide Tiles



2D Variable-Sized Tiles



COO + BCOO formats provide higher non-zero element balance across PIM cores + threads than CSR + BCSR, respectively.

Comparison of Compressed Formats

2048 PIM Cores, 32-bit integer

1D

2D Equally-Sized

Key Takeaway 3

Speedup

The **compressed matrix format** used to store the input matrix **determines** the **data partitioning** across DRAM banks of PIM-enabled memory. As a result, it affects the **load-balance** across PIM cores (and threads of a PIM core) with corresponding **performance** implications.

regular matrices

scale-free
matrices

regular matrices

scale-free
matrices

2D Equally-Wide

2D Variable-Sized

Recommendation 3

Speedup

Design **compressed** data structures that can be **effectively** partitioned across DRAM banks, with the goal of providing **high computation balance** across PIM cores (and threads of a PIM core).

regular matrices

scale-free
matrices

regular matrices

scale-free
matrices

End-to-End Performance

1

Load the
input vector

2

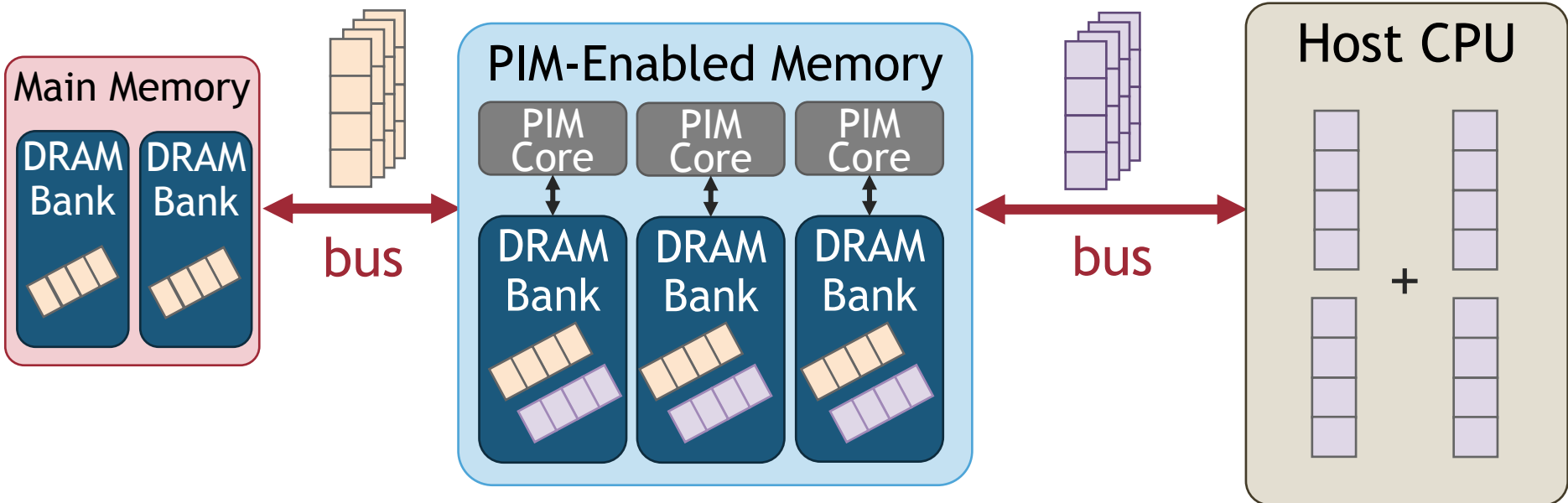
Execute the
kernel

3

Retrieve the
partial results

4

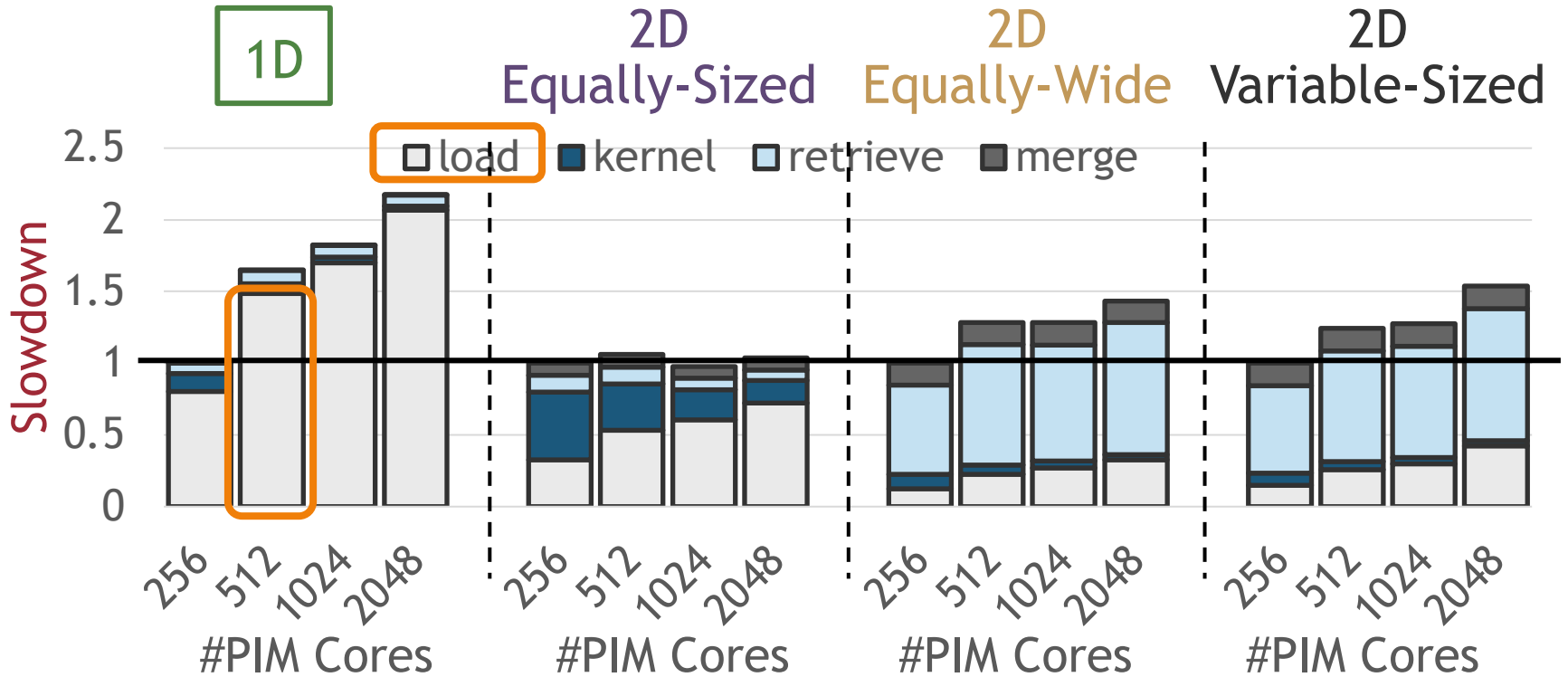
Merge the
partial results



Scalability

COO format, 32-bit integer

The scalability is limited by the **load** time



1D: #bytes to **load** the input vector grows **linearly** to #PIM cores

Scalability

COO format, 32-bit integer

Key Takeaway 4

The 1D-partitioned kernels are severely **bottlenecked** by the high data transfer costs to **broadcast** the whole **input** vector **into DRAM banks** of all PIM cores, through the narrow off-chip memory bus.



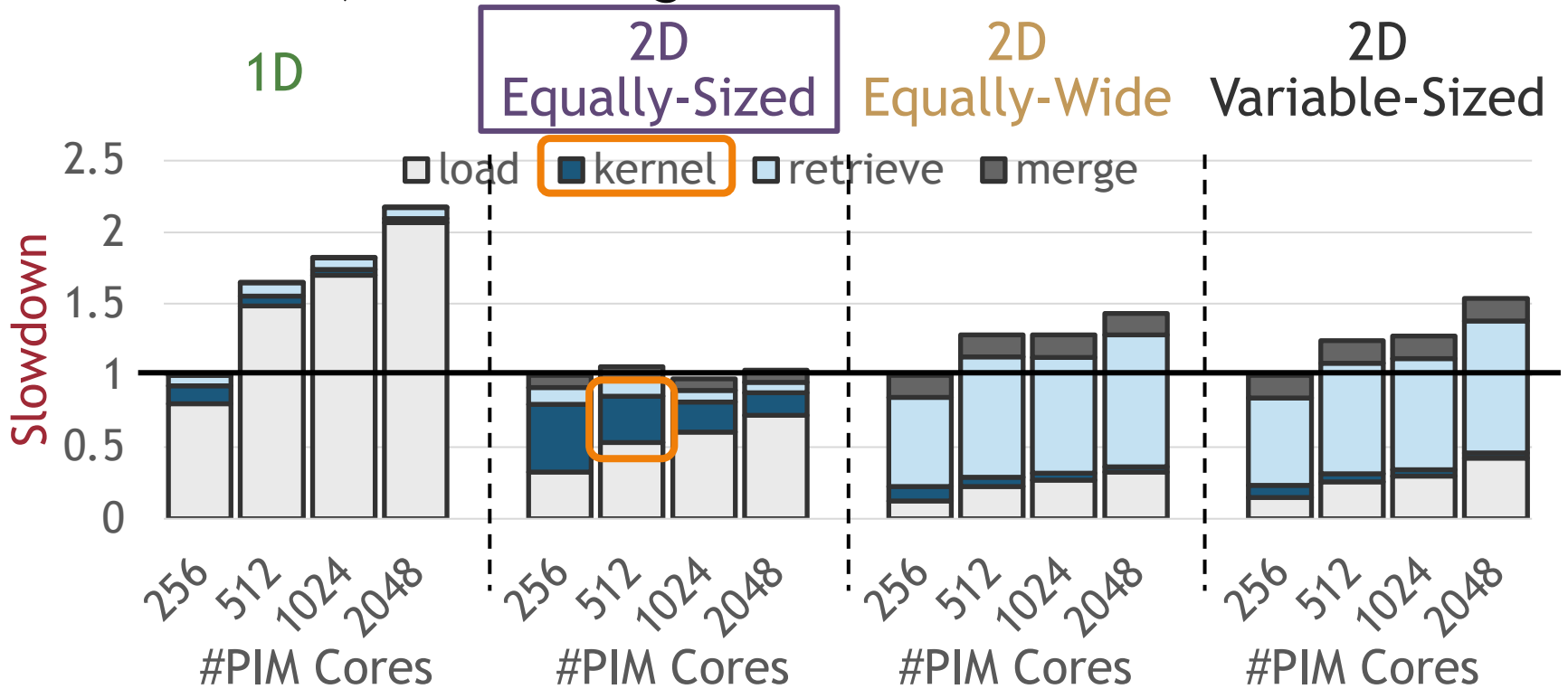
Recommendation 4

Optimize the **broadcast collective** collective in data transfers to PIM-enabled memory to efficiently copy the **input data** into DRAM banks in the PIM system.

Scalability

COO format, 32-bit integer

The scalability is limited by the **kernel** time

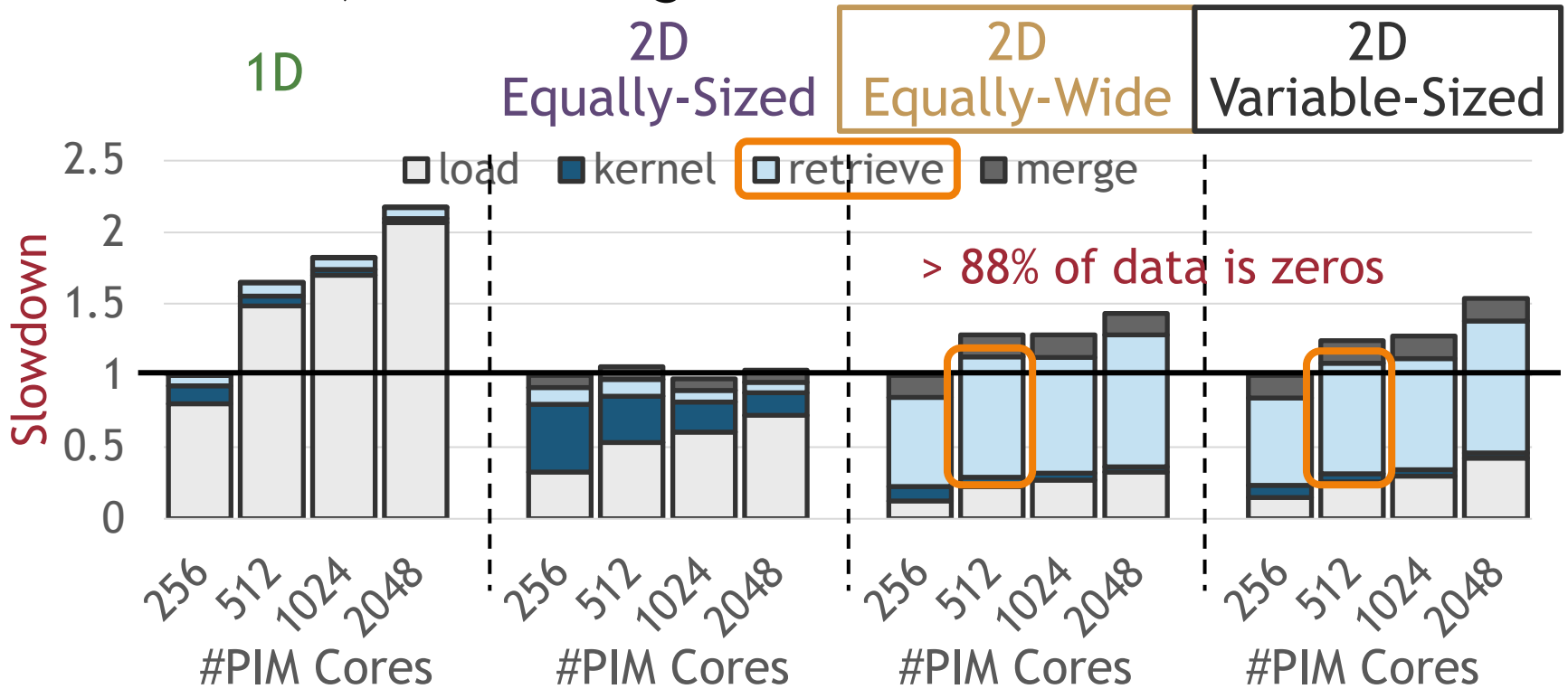


2D Equally-Sized: **kernel** time is limited by only a **few** PIM cores assigned to the 2D tiles with the **largest #NNZs**

Scalability

COO format, 32-bit integer

The scalability is limited by the **retrieve** time



2D Equally-Wide + 2D Variable-Sized:

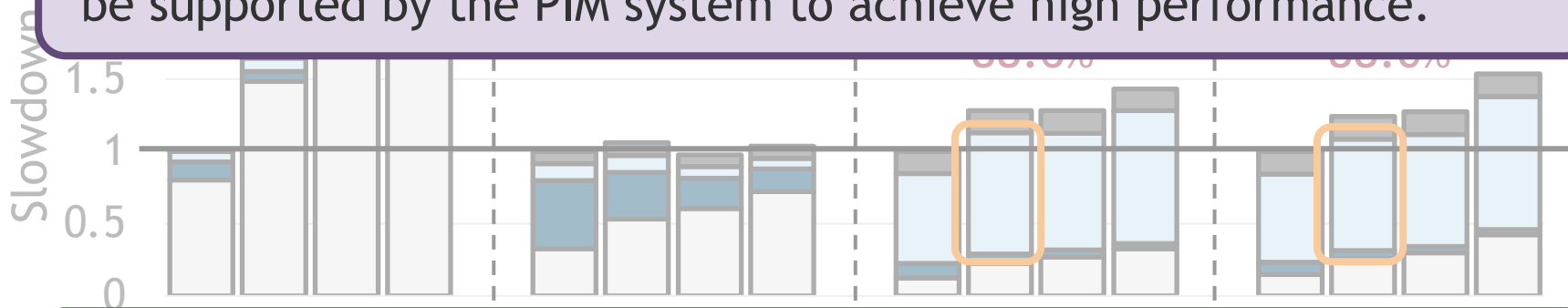
high amount of **zero padding** to **gather** the output vector → **parallel** transfers supported at **rank granularity** = 64 PIM cores

Scalability

COO format, 32-bit integer

Key Takeaway 5

The 2D equally-wide and variable-sized kernels need **fine-grained parallel data transfers** at DRAM bank granularity (**zero padding**) to be supported by the PIM system to achieve high performance.

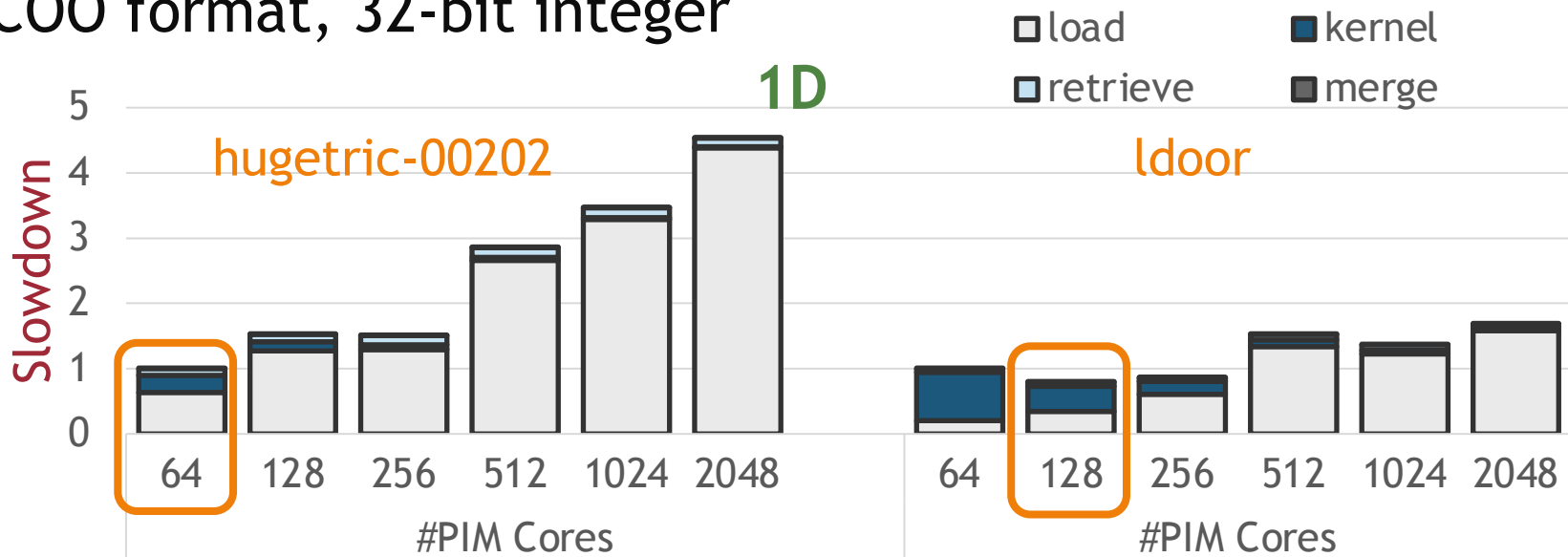


Recommendation 5

Optimize the **gather collective** operation at **DRAM bank granularity** in data transfers from PIM-enabled memory to efficiently retrieve the **output results** to the host CPU.

Comparison of Sparse Matrices

COO format, 32-bit integer



Best-performing = 64 PIM cores

Best-performing = 128 PIM cores

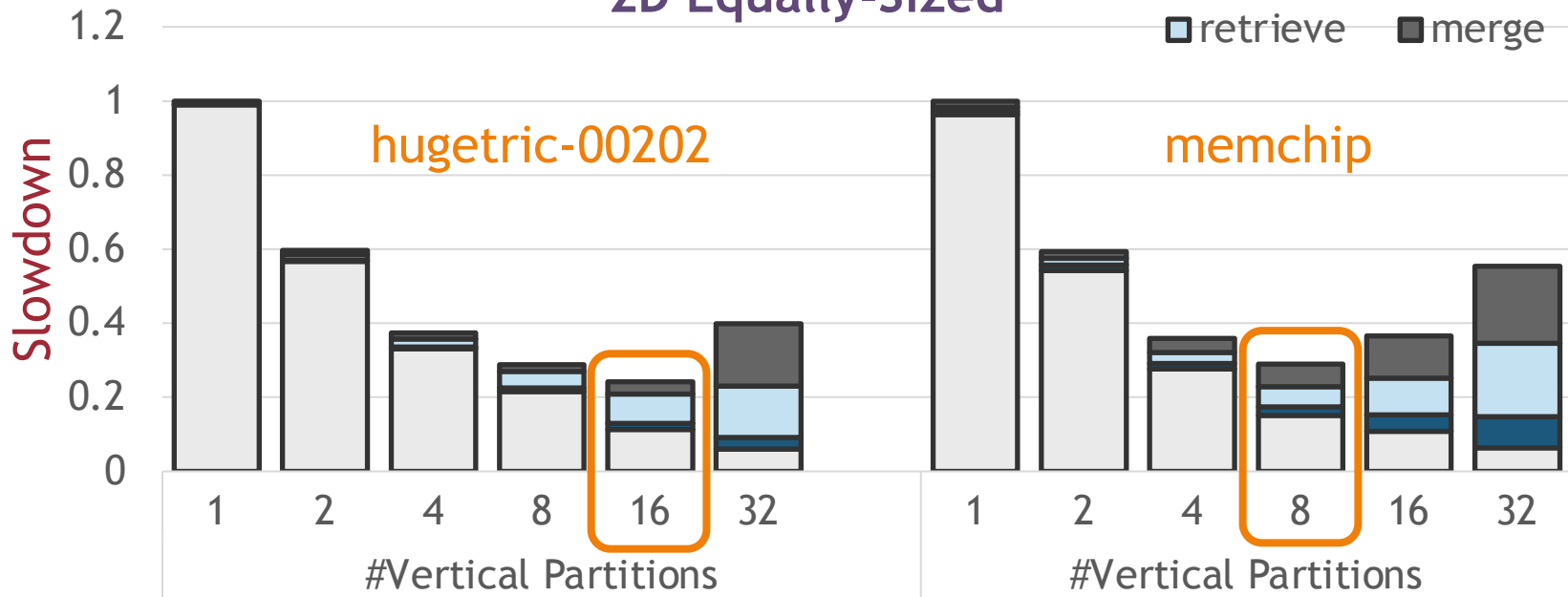
1D: #PIM cores that provides the best performance depends on the sparsity pattern of the input matrix

Comparison of Sparse Matrices

2048 PIM cores, COO format, 32-bit integer

2D Equally-Sized

load kernel
retrieve merge



Best-performing = 16 vertical part.

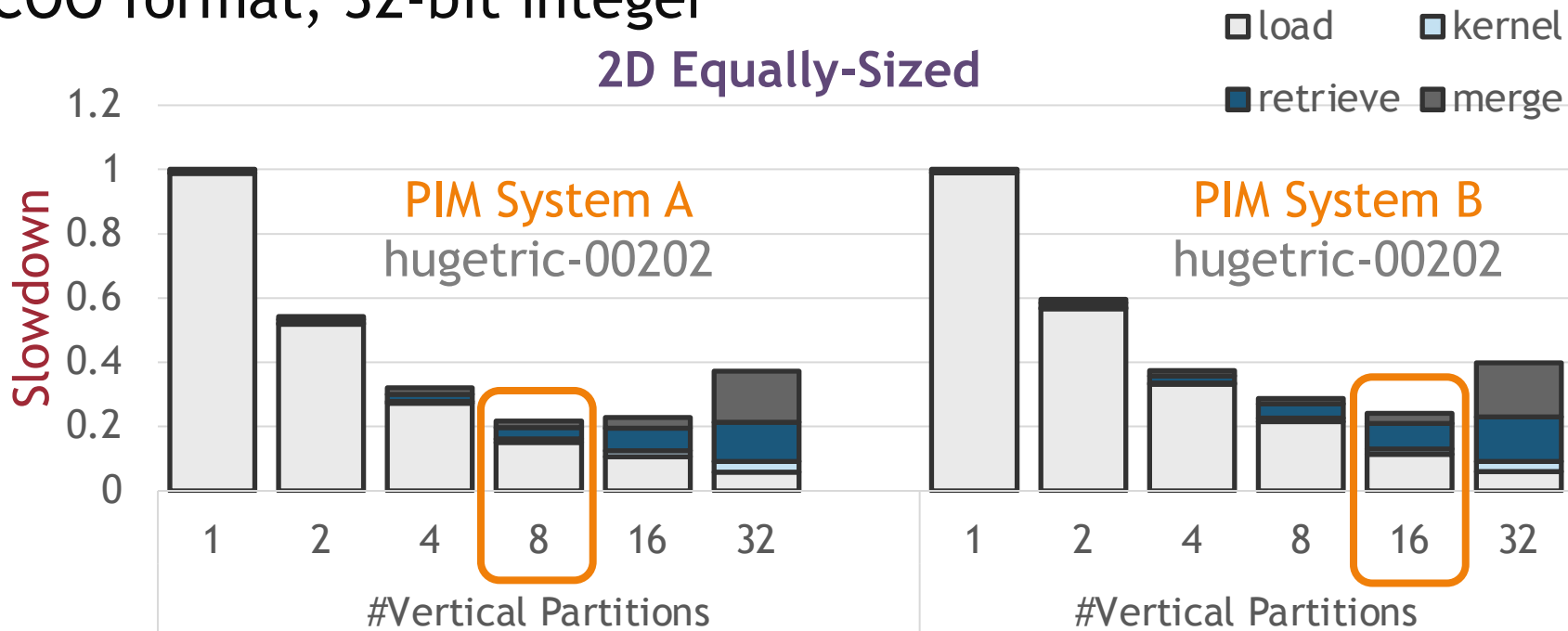
Best-performing = 8 vertical part.

2D: #vertical partitions that provides the best performance depends on the sparsity pattern of the input matrix

Comparison of PIM Systems

COO format, 32-bit integer

2D Equally-Sized



Best-performing = 8 vertical part.

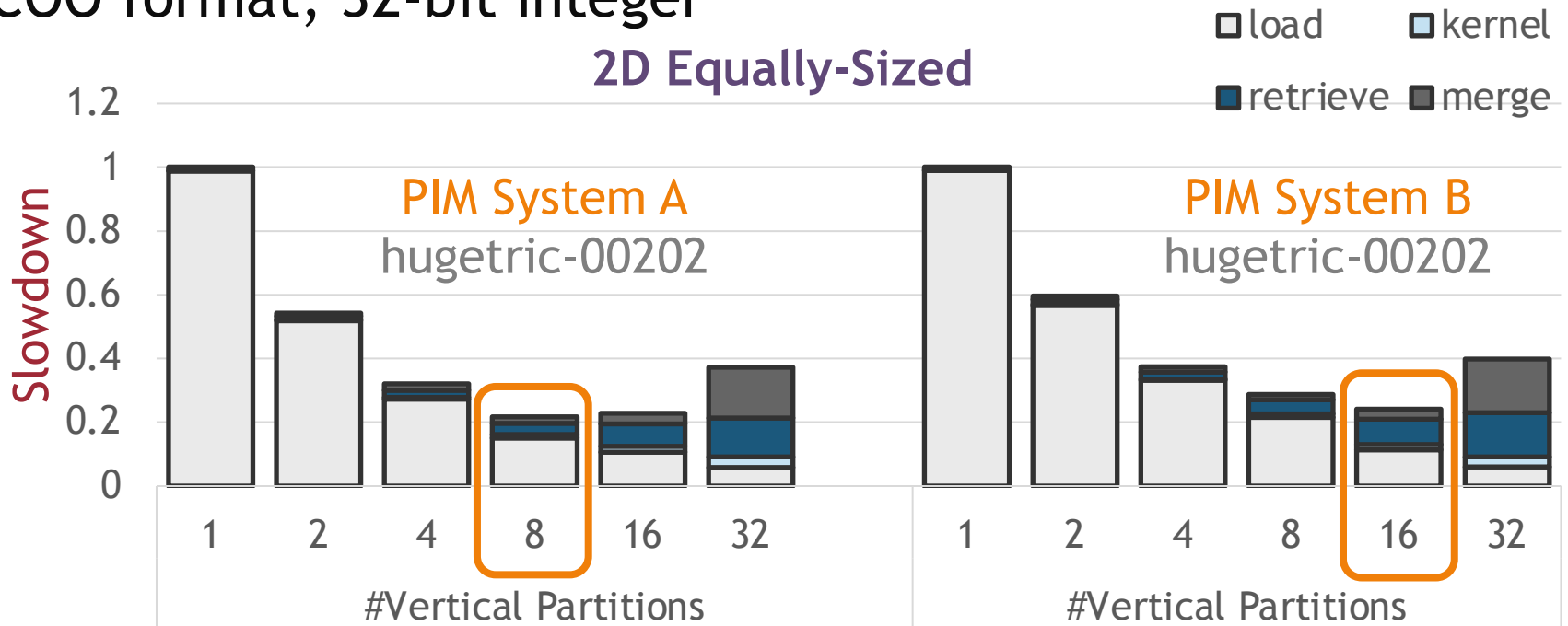
Best-performing = 16 vertical part.

System	PIM Cores	PIM Band.	Host CPU	Bus Band.
PIM A	2048 @350 MHz	1.43 TB/s	Intel Xeon Silver 4110 @2.1 GHz	23.1 GB/s
PIM B	2048 @425 MHz	1.78 TB/s	Intel Xeon Silver 4215 @2.5 GHz	21.8 GB/s

Comparison of PIM Systems

COO format, 32-bit integer

2D Equally-Sized



Best-performing = 8 vertical part.

Best-performing = 16 vertical part.

2D: #vertical partitions that provides the best performance depends on the underlying hardware characteristics

Various Matrices and PIM Systems

COO format, 32-bit integer

load kernel
retrieve merge

1D

5

Key Takeaway 6

There is **no one-size-fits-all** parallelization approach for SpMV, since the performance of each scheme **depends** on the characteristics of the **input matrix** and the underlying **PIM hardware**.

#PIM Cores

#PIM Cores

2D Equally-Sized

2D Equally-Sized

load kernel retrieve merge

load kernel retrieve merge

1.2

1.2

PIM System A

PIM System B

Recommendation 6

Design **adaptive** algorithm that **tune** their configuration to the **particular patterns** of each input given and the **characteristics** of the PIM hardware.

Slowdown

1 2 4 8 16 32

1 2 4 8 16 32

1 2 4 8 16 32

1 2 4 8 16 32

hugetric-0020

memchip

hugetric-0020

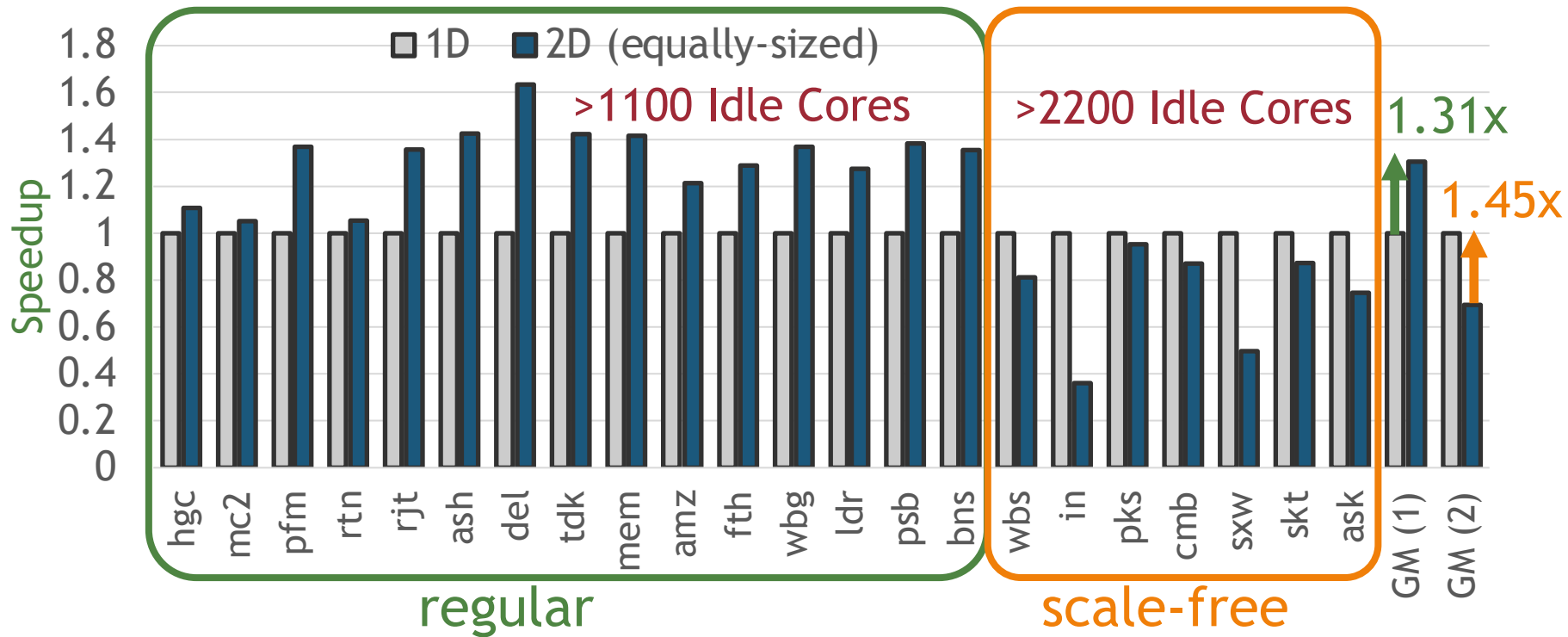
memchip

#Vertical Partitions

#Vertical Partitions

1D vs 2D

Up to 2528 PIM Cores, 32-bit float

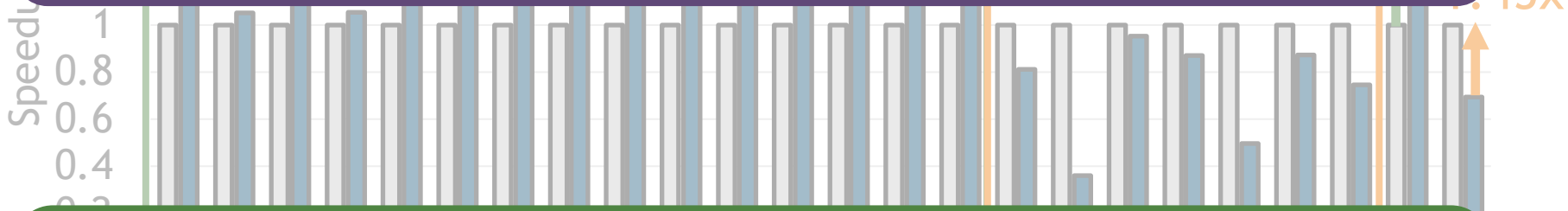


Best-performing SpMV execution:
trades off computation with lower data transfer costs

1D vs 2D

Key Takeaway 7

Expensive **data transfers** to/from PIM-enabled memory performed via the narrow memory bus impose significant **performance overhead** to end-to-end SpMV execution. Thus, it is hard to **fully exploit** all available PIM cores of the system.

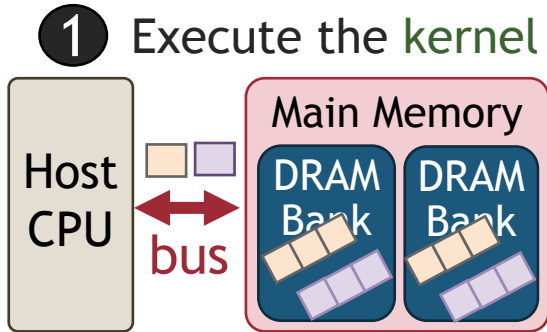


Recommendation 7

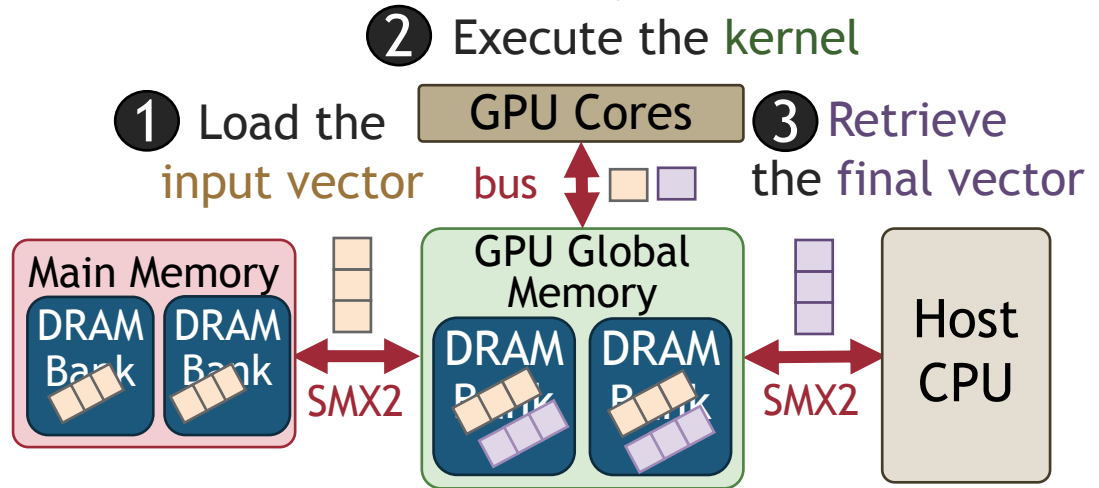
Design **high-speed communication channels** and **optimized libraries** in data transfers to/from PIM-enabled memory, provide **hardware support** to effectively **overlap** computation with data transfers in the PIM system, and/or **integrate** PIM-enabled memory as the main **memory** of the system.

SpMV Execution on Various Systems

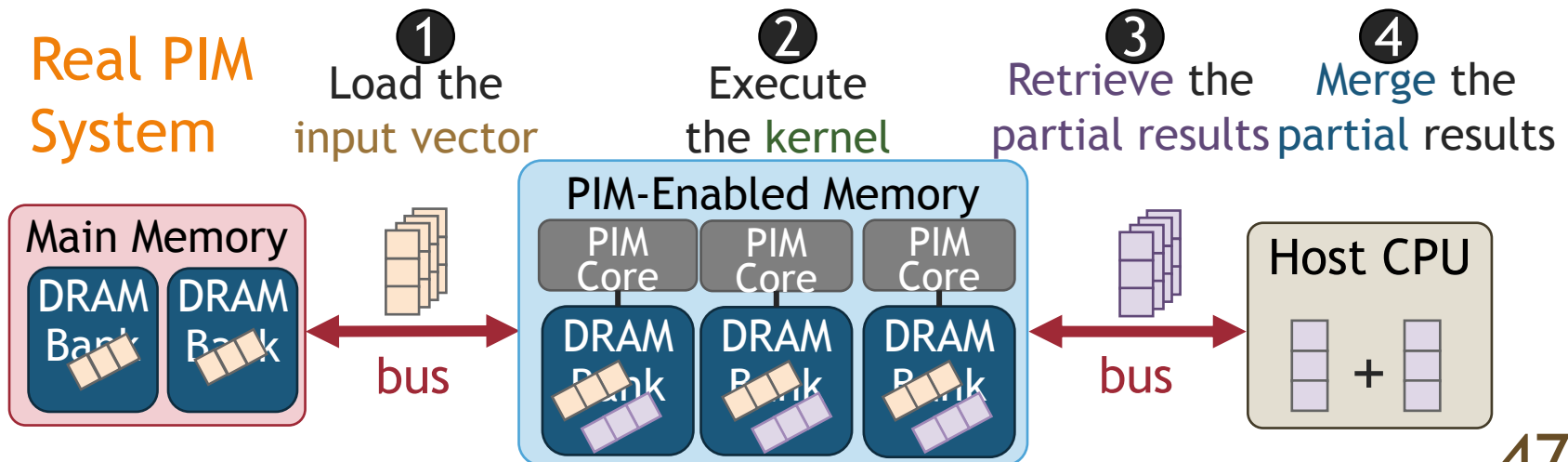
CPU System



GPU System



Real PIM System



CPU/GPU Comparisons

System		Peak Performance	Bandwidth	TDP	
CPU	Intel Xeon Silver 4110	660 GFlops	23.1 GB/s	2x85 W	Processor-Centric
GPU	NVIDIA Tesla V100	14.13 TFlops	897 GB/s	300 W	
PIM	UPMEM 1st Gen.	4.66 GFlops	1.77 TB/s	379 W	Memory-Centric

48

CPU/GPU Comparisons

- **Kernel-Only (COO, 32-bit float):**
 - CPU = 0.51% of Peak Perf.
 - GPU = 0.21% of Peak Perf.
 - PIM (1D) = **50.7%** of Peak Perf.

System		Peak Performance	Bandwidth	TDP	
CPU	Intel Xeon Silver 4110	660 GFlops	23.1 GB/s	2x85 W	} Processor-Centric
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CPU/GPU Comparisons

- **Kernel-Only (COO, 32-bit float):**
 - CPU = 0.51% of Peak Perf.
 - GPU = 0.21% of Peak Perf.
 - PIM (1D) = **50.7%** of Peak Perf.

- **End-to-End (COO, 32-bit float):**
 - CPU = **4.08 GFlop/s**
 - GPU = 1.92 GFlop/s
 - PIM (1D) = 0.11 GFlop/s

System		Peak Performance	Bandwidth	TDP
CPU	Intel Xeon Silver 4110	660 GFlops	23.1 GB/s	2x85 W
GPU	NVIDIA Tesla V100	14.13 TFlops	897 GB/s	300 W
PIM	UPMEM 1st Gen.	4.66 GFlops	1.77 TB/s	379 W

Processor-Centric

Memory-Centric

CPU/GPU Comparisons

- **Kernel-Energy** (COO, 32-bit float):
 - CPU = 0.247 J
 - GPU = 0.051 J
 - PIM (1D) = 0.179 J

PIM: 1.38x higher energy efficiency over CPU

System		Peak Performance	Bandwidth	TDP	
CPU	Intel Xeon Silver 4110	660 GFlops	23.1 GB/s	2x85 W	} Processor-Centric
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CPU/GPU Comparisons

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System	Peak Performance	Bandwidth	TDP
Intel Xeon			

Many more results in the full paper:
<https://arxiv.org/pdf/2201.05072.pdf>

1st Gen.

Centric

Outline

SpMV Kernels for Real PIM Systems

Key Takeaways from Our Study

Conclusion

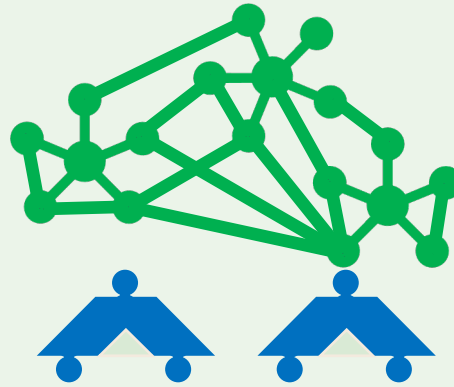
Conclusion

- *SpMV* is a fundamental linear algebra kernel for important applications (HPC, machine learning, graph analytics...)
- *SpMV* is a **highly memory-bound** kernel in processor-centric systems (e.g., CPU and GPU systems)
- Real near-bank PIM systems can tackle the **data movement bottleneck** (high parallelism, large aggregate memory bandwidth)
- Key Contributions:
 - *SparseP*: first **open-source** *SpMV* library for real PIM systems
 - Comprehensive **characterization** and **analysis** of *SpMV* on the first real PIM system
 - **Recommendations** to improve multiple aspects of future PIM hardware and software

Our Work

SparseP: <https://github.com/CMU-SAFARI/SparseP>

Full Paper: <https://arxiv.org/pdf/2201.05072.pdf>



SparseP

Towards Efficient Sparse Matrix Vector Multiplication
on Real Processing-In-Memory Architectures

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